### Numerische Mathematik I

### Homework Problem Sheet 8

For some problems, parts of the solution are already given. Fill in the gaps and complete the proofs where you see a red band at the left margin.

**Introduction.** Quadrature, Tschebyscheff polynomials.

#### **Problem 8.1** Error of Quadrature

Let  $[a, b] \subset \mathbb{R}$  with a < b be a bounded interval and let  $f \in \mathcal{C}^2([a, b])$ .

(8.1a) Show that there is a constant C > 0 independent of a, b and f, such that

$$\left| \int_{a}^{b} f(x) \, \mathrm{d}x - Q_{[a,b]}[f] \right| \leqslant C \cdot |b - a|^{3} \cdot ||f''||_{\mathcal{C}^{0}([a,b])},$$

where  $Q_{[a,b]}[\cdot]$  denotes the rectangle method.

**Solution:** There are many different solutions, we present some of them here:

#### Using polynomial interpolation.

The rectangle method is symmetric, so polynomials of degree smaller than or equal to one are integrated exactly. Let p interpolate f at points  $x_0 = \frac{a+b}{2}$  and  $x_1 \in [a,b] \setminus \{x_0\}$ . Then we have  $\|\omega_2\|_{\mathcal{C}^0([a,b])} \leqslant |b-a|^2$  and we know that  $Q_{[a,b]}[f] = \int_a^b p(x) \,\mathrm{d}x$ .

By [NMI, Thm. 3.6], it follows that  $|f(x)-p(x)| \leq \frac{1}{2}|b-a|^2\|f''\|_{\mathcal{C}^0([a,b])}$ . Estimating the absolute value of the integral then gives

**Using Newton.** Set  $x_0 = \frac{1}{2}(a+b)$ . Then by the theorem about quadrature error of the Midpoint rule in [NMI, Sec. 4.1.2], we have

**Using Taylor.** Set again  $x_0 = \frac{1}{2}(a+b)$ . We have  $f(x) = f(x_0) + hf'(x_0) + \frac{1}{2}(x-x_0)^2 f''(\xi(x))$  hence

**Using Gauss.** The trapezoidal rule is exactly the Gaussian quadrature for n=0. By [NMI, Thm. 4.18], we have

$$\left| \int_{-1}^{1} \widetilde{f}(t) \, dt - 2\widetilde{f}(0) \right| \leq \frac{1}{3} \|\widetilde{f}''\|_{\mathcal{C}^{0}([-1,1])}.$$

Set  $x_0 = \frac{1}{2}(a+b)$  again. With  $\widetilde{f}(t) := f(x_0 + (b-x_0)t)$  we get

$$\frac{b-a}{2} \left( 2f(x_0) - \int_a^b f(x) \, dx \right) = 2\tilde{f}(0) - \int_{-1}^1 \tilde{f}(t) \, dt$$

and

$$\|\widetilde{f}''\|_{\mathcal{C}^0([-1,1])} = \frac{(b-a)^2}{2^2} \|f''\|_{\mathcal{C}^0([a,b])},$$

so the result follows.

**(8.1b)** For  $f \in C^0([0,1])$  and  $h = (N-1)^{-1}$ , let  $T_h[f]$  denote the iterated trapezoidal rule on N equidistant sampling points in the interval [0,1].

Show that for  $f(x) := x^{\alpha}$ , where  $0 < \alpha < 1$ , we have

$$\left| \int_0^1 f(x) \, \mathrm{d}x - T_h[f] \right| = \mathcal{O}(h^{\alpha+1}) \quad \text{for } h \to 0.$$

What changes if  $f(x) = x^{\alpha}g(x)$ , where  $g \in \mathcal{C}^3([0,1])$ ?

**Solution:** Let  $T_J^{(K)}[f]$  denote the trapezoidal rule with K equidistant sampling points on the interval J.

By [NMI, Thm. 4.4] with h = 1/(N-1), we get that the trapezoidal rule satisfies

$$\left| \int_{a}^{b} f(x) \, \mathrm{d}x - T_{[a,b]}^{(N)}[f] \right| \leqslant \frac{h^{2}}{12} (b-a) \max_{x \in [a,b]} |f''(x)|$$

for  $f \in \mathcal{C}^2([a,b])$  – but this is the crucial point:  $f \notin \mathcal{C}^2([a,b])$ , so this theorem cannot be applied. Even just applying it to an interval [h,1] does not help because  $f''(h) \propto h^{\alpha-2}$ .

So let  $[a_i, b_i] := h[i-1, i]$  for i = 1, ..., N-1 be the subintervals of the trapezoidal rule on N equidistant sampling points. On each of the N subintervals  $J = [a_i, b_i] = [b_i - h, b_i]$ , i = 2, ..., N-1, the trapezoidal rule satisfies

For  $f(t)=t^{\alpha}$  with  $0<\alpha<1$ , we have  $f'(t)=\alpha t^{\alpha-1}$  and  $f''(t)=\alpha(\alpha-1)t^{\alpha-2}$ . Hence f'' is a strictly increasing function on  $(0,\infty)$  and so we have  $f''(\xi)\leqslant Cb_i^{\alpha-2}$  for all  $\xi\in J$ .

Summing this over all N-1 intervals gives

### **Problem 8.2** Order of Quadrature

(8.2a) Calculate the points  $x_0, x_1 \in [-1, 1]$  and the weights  $A, B \in \mathbb{R}$  for the quadrature rule

$$\int_{-1}^{1} f(x) dx \approx Af(x_0) + Bf(x_1)$$

such that the formula has the highest possible degree. What is that degree?

**Solution:** A quadrature formula on n+1 points is of maximum degree 2n+2, in our case (n=1) we get degree  $\leq 4$ . The degree equals 4 if all monomials  $1, x, x^2, x^3$  can be integrated exactly by our formula. This is the case if and only if

$$2 = A + B \tag{8.2.1}$$

$$0 = Ax_0 + Bx_1 (8.2.2)$$

$$\frac{2}{3} = Ax_0^2 + Bx_1^2 \tag{8.2.3}$$

$$0 = Ax_0^3 + Bx_1^3 (8.2.4)$$

From (8.2.2) and (8.2.4) we conclude  $Bx_1(x_0^2-x_1^2)=0$ . Hence one of the following cases must hold:

**(8.2b)** Calculate the *exact* result of  $I = \int_1^5 \left| \frac{1}{2}x - \frac{3}{2} \right|^3 dx$  by using the result we obtained in subproblem (8.2a) and by modifying I in such a way that the integrand becomes a polynomial.

**Solution:** 

From our quadrature formula we get:

#### **Problem 8.3** Iterated Quadrature Formulas

Let  $x_i$ , i = 0, ..., n, with  $-\infty < x_0 < x_1 < ... < x_n < \infty$  be fixed real numbers in arithmetic progression and let f be a smooth function on the interval  $[x_0, x_n]$ .

**(8.3a)** Check that  $T_{h/2} = \frac{1}{2}(T_h + M_h)$ , where

$$T_h[f] = \frac{h}{2} \left( f(x_0) + 2 \sum_{i=1}^{n-1} f(x_i) + f(x_n) \right)$$
 and  $M_h[f] = h \sum_{i=0}^{n-1} f(x_i + h/2)$ 

are the iterated trapezoidal rule and the iterated rectangle method, respectively.

**Solution:** We have

$$T_h[f] = \frac{h}{2} \left( f(x_0) + 2 \sum_{i=1}^{n-1} f(x_i) + f(x_n) \right)$$

$$M_h[f] = h \sum_{i=0}^{n-1} f(x_{i+1/2}) \quad \text{where} \quad x_{i+1/2} := \frac{x_i + x_{i+1}}{2}.$$

Hence

**(8.3b)** Explicitly calculate one step of Romberg's method with values  $T_h[f]$  and  $T_{h/2}[f]$ . Rearrange the result to a quadrature formula of step width h.

**Solution:** For f sufficiently smooth, we get an asymptotic expansion of the form

$$T_h[f] = I + c_2 h^2 + \mathcal{O}(h^4)$$
 and  $T_{h/2}[f] = I + c_2 \frac{h^2}{4} + \mathcal{O}(h^4)$ ,

where I is the exact value of the integral. We now combine both terms linearly such that the terms with order  $h^4$  cancel out. Set  $T_{00} = T_h[f]$  and  $T_{10} = T_{h/2}[f]$ . Then we get

**(8.3c)** What is the asymptotic behaviour of the error of the new method?

**Solution:** 

# **Problem 8.4 Chebychev Polynomials**

In this problem you will meet a special set of polynomials that form a basis for the spaces  $\mathbb{P}_{n-1}$  of polynomials of degree  $< n, n \in \mathbb{N}$ . With respect to an also special set of interpolation nodes this basis has rather desirable properties.

Throughout this problem, for  $n \in \mathbb{N}_0$  we set

$$T_n(t) := \cos(n \arccos(t)), \quad -1 \le t \le 1,$$

and for  $n \in \mathbb{N}$ 

$$\mathcal{Z}_n := \left\{ x_k := \cos\left( (k + \frac{1}{2}) \frac{\pi}{n} \right), \ k = 0, \dots, n - 1 \right\}.$$
 (8.4.1)

- **(8.4a)** Read through section 3.8 of the lecture notes.
- **(8.4b)** Show that  $T_0(t) = 1$ ,  $T_1(t) = t$ , and

$$T_{n+1}(t) + T_{n-1}(t) = 2tT_n(t) , \quad n \ge 1 , \quad -1 \le t \le 1 .$$
 (8.4.2)

(8.4c) Show that for  $n \ge 1$  the derivatives of the functions  $T_n$  satisfy the recursion

$$2T_n(t) = \frac{1}{n+1} \frac{d}{dt} T_{n+1}(t) - \frac{1}{n-1} \frac{d}{dt} T_{n-1}(t) .$$

- (8.4d) Show that  $T_n$  for  $n \ge 1$ , is a polynomial of degree n with leading coefficient  $2^{n-1}$ .
- **(8.4e)** Show that

$$\mathcal{Z}_n = \{ t \in \mathbb{R} : T_n(t) = 0 \} .$$

(8.4f) Write a MATLAB function

function 
$$y = evalT(n,t)$$

that computes  $y_i := T_n(t_i)$ ,  $n \in \mathbb{N}_0$ , for arguments  $t_i$ , i = 1, ..., m, passed in the row vector t. No special functions like  $\cos$  and its inverse must be used.

The results are to be returned in the row vector y. What is the asymptotic computational effort of evalT for  $n \to \infty$  and  $m \to \infty$ ?

HINT: Use the recursion formula (8.4.2) for  $T_n$ .

(8.4g) We consider the interpolation problem with  $V = \mathbb{P}_{n-1}$  and set of interpolation points  $\mathcal{Z}_n$ .

Let  $\mathbf{A} \in \mathbb{R}^{n \times n}$  be the interpolation matrix for the basis  $\mathcal{B} = \{T_0, T_1, \dots, T_{n-1}\}$  of V (i.e. the (j+1)th column of  $\mathbf{A}$  equals  $(T_j(x_0), \dots, T_j(x_{n-1}))^{\top}$ . Show that there is a regular diagonal matrix  $\mathbf{D}$  such that  $\mathbf{A}\mathbf{D}$  is an orthogonal matrix.

**(8.4h)** Using the result of subproblem (8.4g), write an efficient (in terms of computational effort and memory!) MATLAB function

function 
$$c = chebcoeff(y)$$

that computes the coefficients  $c_j$ ,  $j = 0, \dots, n-1$ , in the representation

$$p(t) = \sum_{j=0}^{n-1} c_j T_j(t)$$

of the polynomial interpolant  $p \in \mathbb{P}_{n-1}$  through the points  $(x_i, y_i)$ ,  $i = 0, \ldots, n-1$ , where the nodes  $x_i$  are defined in (8.4.1). The data  $y_i$  are passed in the row vector y, and the coefficients are returned in the row vector z.

(8.4i) Based on (8.4.2) write a MATLAB function

function 
$$y = chebsum(c, t)$$

that computes

$$y_i = \sum_{j=0}^{n-1} c_j T_j(t_i)$$

for i = 1, ..., m. The values  $t_i$  are the components of the row vector t, and the values  $y_i$  are made available in the row vector y.

What is the asymptotic computational effort of chebsum for  $n \to \infty$  and  $m \to \infty$ ?

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MATLAB: Submit all files in the online system. Include the files that generate the plots. Label all your plots. Include commands to run your functions. Comment on your results.

## References

[NMI] Lecture Notes for the course "Numerische Mathematik I".

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