

The knitting algorithm

Finite root systems

Let V be an Euclidean space with inner product (\cdot, \cdot) .

Definition 1. A subset R of V is called a **finite root system**, if it satisfies for $\alpha, \beta \in R$:

$$(R1) \quad R \text{ is finite, } \langle R \rangle = V, 0 \notin R$$

$$(R2) \quad \lambda \alpha \in R, \lambda \in \mathbb{R} \Rightarrow \lambda = \pm 1$$

$$(R3) \quad \forall \gamma \in R : \sigma_\alpha(\gamma) \in R, \text{ where } \sigma_\alpha(\gamma) := \gamma - \frac{2(\gamma, \alpha)}{(\alpha, \alpha)}.$$

$$(R4) \quad \langle \alpha, \beta \rangle := \frac{2(\beta, \alpha)}{(\alpha, \alpha)} \in \mathbb{Z}$$

When we talk of a root system in this paper, we actually mean a finite root system. The **rank** of a root system R is the dimension of V .

Definition 2. R is called **irreducible** if there is no representation $R = R_1 \cup R_2$ with two proper subsets R_1 and R_2 and $\forall x \in R_1 \forall y \in R_2 : (x, y) = 0$.

Definition 3. A subset Δ of a root system R which is a basis for V is called **basis** of R , if for all $\beta = \sum_{\alpha \in \Delta} k_\alpha \alpha$ the coefficients k_α are integers and all less or equal to 0 or all greater or equal to 0. The elements in the basis Δ are called **simple roots**. Because Δ is a basis of V , the representation $\beta = \sum_{\alpha \in \Delta} k_\alpha \alpha$ is well-defined. If $\sum_{\alpha \in \Delta} k_\alpha > 0$ ($\sum_{\alpha \in \Delta} k_\alpha < 0$), then the root is called **positive (negative) root**.

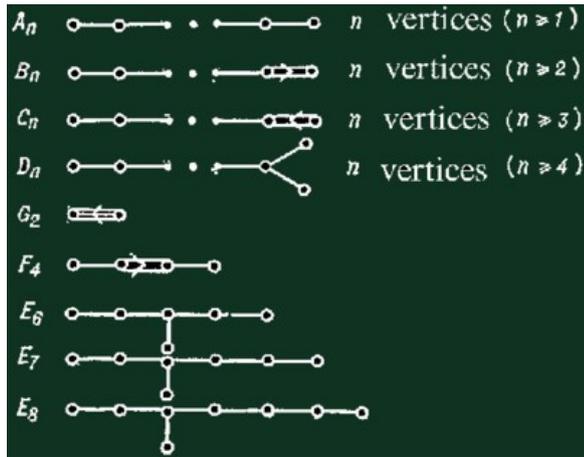
Proposition 1. Every finite root system has a basis.

Dynkin-diagram Let R be a finite root system.

Definition 4. Let $\alpha_1, \dots, \alpha_l$ be simple roots of R of rank l . Then the $l \times l$ -matrix $(\langle \alpha_i, \alpha_j \rangle)_{ij}$ is called the **Cartan-matrix** of R .

Proposition 2. The Cartan-matrix of R is well-defined, i.e. independent of the choice of a basis of R .

Definition 5. Let $\alpha_1, \dots, \alpha_l$ be the simple roots of R . We can construct a graph as follows: Take l vertices. Between vertex i and j we draw $\langle \alpha_i, \alpha_j \rangle \langle \alpha_j, \alpha_i \rangle$ edges (well-defined, because Cartan-matrix is independent of basis). There is an arrow on the edges from the larger root to the shorter one. This graph is called **Dynkin-diagram** of R .



Remark 1. R is irreducible if and only if the Dynkin-diagram of R is connected.

Theorem 1. Let R be an irreducible root system. Then its Dynkin-diagram is one of the following (and all diagrams can be realized as diagrams of irreducible root systems):

Remark 2. These diagrams also classify the finite dimensional simple complex Lie-Algebras.

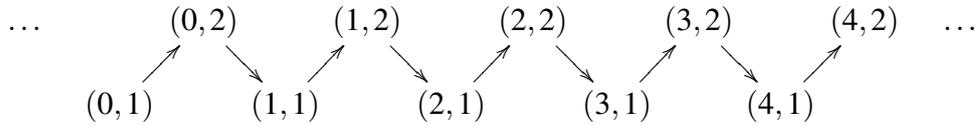
Definition 6. Let Λ be a commutative ring. An **algebra** over Λ is a Λ -module A together with a bilinear map $A \times A \rightarrow A$.

Definition 7. A **cluster algebra** is a commutative \mathbb{Q} -algebra endowed with a family of generators, the so-called **cluster variables**, grouped into overlapping subsets of fixed finite cardinality, which are constructed recursively using mutations. If the set of cluster variables is finite, we speak of **cluster-finite cluster algebras**.

Theorem 2. The cluster-finite cluster algebras are parametrized by the finite root systems, i.e. by the Dynkin-diagrams above.

The algorithm

The **knitting algorithm** shows how to get the cluster-finite cluster algebra \mathcal{A}_Δ corresponding to a Dynkin-diagram Δ . Let's start with the simplest non trivial Dynkin-diagram. $A_2 : 1 \longrightarrow 2$. Now, we construct the so-called **repetition** $\mathbb{Z}A_2$. First, we draw the product $\mathbb{Z} \times A_2$. Then we draw for each arrow $(n, i) \longrightarrow (n, j)$ a new one $(n, j) \longrightarrow (n + 1, i)$ and we get the repetition.



Now, we assign to each vertex of the repetition a cluster variable. We start with x_1 and x_2 at the 0-th copy of A_2 .



We can construct x'_1, x'_2, x''_1, \dots by "knitting" from the left to the right. To compute x'_1 , we divide the sum of the immediate predecessor of x'_1 and 1 by the left translate of x'_1 . That gives us: $x'_1 = \frac{x_2+1}{x_1}$. We do the same procedure for x'_2, x''_1, \dots . So, finally we get: $x'_2 = \frac{x_1+x_2+1}{x_1x_2}$, $x''_1 = \frac{x_1+1}{x_2}$, $x''_2 = x_1$ and $x'''_1 = x_2$. Obviously, from now on the whole thing will repeat. Thus, we have 5 cluster variables x_1, x_2, x'_1, x'_2 and x''_1 . So, the cluster algebra \mathcal{A}_{A_2} is a \mathbb{Q} -subalgebra of $\mathbb{Q}(x_1, x_2)$ generated by the cluster variables above. We can see that all denominators of the cluster variables are monomials. This can be generalized:

- (1) All denominators of all cluster variables are monomials. This holds for all cluster algebras. It's called **Laurent phenomenon**.
- (2) The computation is **periodic** and we always get only finitely many cluster variables (as expected). Furthermore, the algorithm can easily be extended from Dynkin-diagrams to arbitrary trees, and then one can characterize the Dynkin diagrams by the periodicity of the algorithm.
- (3) We have $5 = 2 + 3$ cluster variables. 2 initial ones x_1, x_2 and 3 non initial ones x'_1, x'_2 and x''_1 . The non initial ones are in natural bijection with the positive roots of the corresponding root system of A_2 . To see that there is a bijection, consider the denominators of the cluster variables.

$$\frac{1}{x_1^{d_1} x_2^{d_2}} \mapsto d_1 \alpha_1 + d_2 \alpha_2$$

Fomin-Zelevinsky showed that this generalizes to arbitrary Dynkin-diagrams. In general, the number of cluster variables in the cluster algebra \mathcal{A}_Δ is equal to the sum of the rank of Δ and the number of positive roots of Δ .

More examples (for $\Delta = A_3$ and $\Delta = G_2$) are in the lecture notes of Bernhard Keller.