

## AN UPPER BOUND OF A GENERALIZED UPPER HAMILTONIAN NUMBER OF A GRAPH

MARTIN DZÚRIK

**ABSTRACT.** In this article we study graphs with ordering of vertices, we define a generalization called a pseudoordering, and for a graph  $H$  we define the  $H$ -Hamiltonian number of a graph  $G$ . We will show that this concept is a generalization of both the Hamiltonian number and the traceable number. We will prove equivalent characteristics of an isomorphism of graphs  $G$  and  $H$  using  $H$ -Hamiltonian number of  $G$ . Furthermore, we will show that for a fixed number of vertices, each path has a maximal upper  $H$ -Hamiltonian number, which is a generalization of the same claim for upper Hamiltonian numbers and upper traceable numbers. Finally we will show that for every connected graph  $H$  only paths have maximal  $H$ -Hamiltonian number.

### 1. INTRODUCTION

In this article we study a part of graph theory based on an ordering of vertices. We define a generalization called a pseudoordering of a graph. We will show how to generalize a Hamiltonian number, for a graph  $H$  we define the  $H$ -Hamiltonian number of a graph  $G$  and we will show that this concept is a generalization of both the Hamiltonian number and the traceable number. We get them by a special choice of graph  $H$ . Furthermore, we will study a maximalization of upper  $H$ -Hamiltonian number for a fixed number of vertices. We will show that, for a fixed number of vertices, each path has a maximal upper  $H$ -Hamiltonian number. From the definition it will be obvious that a lower bound of the  $H$ -Hamiltonian number is the number of edges  $|E(H)|$  and the graph  $G$  has a minimal lower  $H$ -Hamiltonian number if and only if  $H$  is a subgraph of  $G$ . Now we can say that  $G$  having a maximal upper  $H$ -Hamiltonian number is dual to  $H$  being a subgraph of  $G$ . Furthermore, by above for every two finite graphs  $G$  and  $H$  such that  $G$  is connected satisfying  $|V(G)| = |V(H)|$  and  $|E(G)| = |E(H)|$ , we get that  $G \cong H$  if and only if the lower  $H$ -Hamiltonian number of  $G$  is  $|E(H)|$ .

In [2] it is proved that  $G$  has a maximal upper traceable number if and only if  $G$  is a path. The same is proved for Hamiltonian number. We will show that for

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$H$  connected  $G$  has a maximal  $H$ -Hamiltonian number if and only if  $G$  is a path. This shows that this generalization of ordering of vertices is natural.

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In this article we will study a generalization of Hamiltonian spectra of undirected finite graphs. Recall that, a graph  $G$  is a pair

$$G = (V(G), E(G)),$$

where  $V(G)$  is a finite set of vertices of  $G$  and  $E(G) \subseteq V(G) \times V(G)$ , a symmetric Antireflexive relation, is a set of edges. We will denote an edge between  $v$  and  $u$  by  $\{v, u\}$ .

Recall that, an *ordering* on the graph  $G$  is a bijection

$$f: \{1, 2, \dots, |V(G)|\} \rightarrow V(G),$$

we denote

$$s(f, G) = \sum_{i=1}^{|V(G)|} \rho_G(f(i), f(i+1)),$$

$$\bar{s}(f, G) = \sum_{i=1}^{|V(G)|-1} \rho_G(f(i), f(i+1)),$$

where  $\rho_G(x, y)$  is the distance of  $x, y$  in the graph  $G$  and  $f(|V(G)| + 1) := f(1)$ , for better notation. We will write only  $s(f)$ ,  $\bar{s}(f)$  if the graph is clear from context. Then

$$\{s(f, G) \mid f \text{ ordering on } G\}$$

$$\{\bar{s}(f, G) \mid f \text{ ordering on } G\}$$

are the *Hamiltonian spectrum* of the graph  $G$  and the *traceable spectrum* of the graph  $G$ , respectively.

We want to generalize the notion of an ordering of a graph.

**Definition 1.1.** Let  $G, H$  be graphs such that  $|V(G)| = |V(H)|$  and  $f: V(H) \rightarrow V(G)$  is a bijection, then we call  $f$  a *pseudoordering* on the graph  $G$  (by  $H$ ), denote

$$s_H(f, G) = \sum_{\{x, y\} \in E(H)} \rho_G(f(x), f(y)),$$

where  $\rho_G(x, y)$  is the distance of  $x, y$  in the graph  $G$ . We will call  $s_H(f, G)$  the sum of the pseudoordering  $f$ . Then

$$\{s_H(f, G) \mid f \text{ pseudoordering on } G \text{ by } H\}$$

is the  *$H$ -Hamiltonian spectrum* of the graph  $G$ .

The minimum and the maximum of a Hamiltonian spectrum and of a traceable spectrum are called the (*lower*) *Hamiltonian number* and the *upper Hamiltonian number*, respectively. Furthermore, the (*lower*) *traceable number* and the *upper traceable number* of a graph  $G$  are denoted by

$$\begin{aligned}
 h(G) &= \min\{s(f, G) \mid f \text{ ordering on } G\}, \\
 h^+(G) &= \max\{s(f, G) \mid f \text{ ordering on } G\}, \\
 t(G) &= \min\{\bar{s}(f, G) \mid f \text{ ordering on } G\}, \\
 t^+(G) &= \max\{\bar{s}(f, G) \mid f \text{ ordering on } G\}.
 \end{aligned}$$

Now we define generalized versions.

**Definition 1.2.**

$$\begin{aligned}
 h_H(G) &= \min\{s_H(f, G) \mid f \text{ pseudoordering on } G\}, \\
 h_H^+(G) &= \max\{s_H(f, G) \mid f \text{ pseudoordering on } G\}.
 \end{aligned}$$

We will call them the *lower H-Hamiltonian number* and the *upper H-Hamiltonian number* of a graph  $G$ , respectively.

Now take  $H = C_{|V(G)|}$ , where  $C_n$  is the cycle with  $n$  vertices. When we denote the vertices of  $C_{|V(G)|}$  by  $\{1, 2, \dots, |V(G)|\}$  we can see that

$$s(f, G) = s_{C_{|V(G)|}}(f, G).$$

Analogously for  $H = P_{|V(G)|-1}$ , where  $P_{n-1}$  is the path of length  $n - 1$ , we get that

$$\bar{s}(f, G) = s_{P_{|V(G)|-1}}(f, G).$$

**Remark 1.3.** The  $C_{|V(G)|}$ -Hamiltonian spectrum of a graph  $G$  is equal to the Hamiltonian spectrum of  $G$  for  $|V(G)| \geq 3$ , and the  $P_{|V(G)|-1}$ -Hamiltonian spectrum of  $G$  is equal to the traceable spectrum of  $G$  for  $|V(G)| \geq 2$ .

**Lemma 1.4.** *Let  $G$  be a connected finite graph and  $H$  be a graph such that  $|V(G)| = |V(H)|$ , then  $h_H(G) = |E(H)|$  if and only if  $H$  is isomorphic to some subgraph of  $G$ .*

**Proof.** Let  $f: V(H) \rightarrow V(G)$  be a pseudoordering satisfying  $s(f, G) = |E(H)|$ , then  $f$  is an injective graph homomorphism. The opposite implication is obvious.  $\square$

**Lemma 1.5.** *Let  $G$  be a connected finite graph and  $H$  be a graph such that  $|V(G)| = |V(H)|$  and  $|E(G)| = |E(H)|$ , then  $h_H(G) = |E(H)|$  if and only if  $H$  is isomorphic to the graph  $G$ .*

**Proof.** The graph  $H$  is isomorphic to a subgraph of  $G$  and furthermore  $|V(G)| = |V(H)|$ ,  $|E(G)| = |E(H)|$ , hence  $H \cong G$ . The opposite implication is obvious.  $\square$

2. MAXIMALIZATION OF THE UPPER  $H$ -HAMILTONIAN NUMBER OF A GRAPH  $G$

In this section we will prove that for every pair of connected graphs  $H, G$  and each pseudoordering  $f$  there exists a pseudoordering

$$g: V(H) \rightarrow \{1, 2, \dots, |V(G)|\}$$

such that

$$s_H(f, G) \leq s_H(g, P_{|V(G)|-1}).$$

At first, let  $G$  be a tree. We will only work with graphs which have at least 2 vertices.

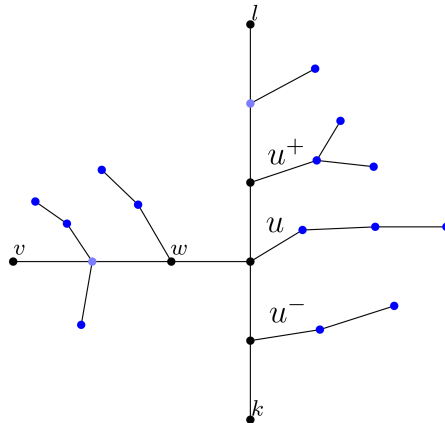
**Definition 2.1.** Let  $G$  and  $H$  be graphs such that  $G$  is connected,  $|V(G)| = |V(H)|$  and  $f: V(H) \rightarrow V(G)$  is a pseudoordering. Furthermore, let  $a, b \in V(G)$ , we define  $a \sim_{H,f} b$  if and only if  $\{f^{-1}(a), f^{-1}(b)\} \in E(H)$ .

**Definition 2.2.** Let  $G$  be a tree such that  $G$  is not a path. Denote three pairwise distinct leaves by  $l, k, v \in V(G)$ . Because  $G$  is not a path then  $G$  has at least 3 leaves, connect  $l, k$  with a path  $l = x_1, x_2, \dots, x_m = k$ . Connect  $v, l$  with a path  $v = y_1, y_2, \dots, y_s = l$  and take the minimum of a set

$$i_m = \min\{i \mid \exists j \in \{1, \dots, m\}, y_i = x_j\}.$$

Take  $j_m$  such that  $y_{i_m} = x_{j_m}$ . Now we define  $u = y_{i_m}, w = y_{i_m-1}, u^+ = x_{j_m-1}, u^- = x_{j_m+1}$ .

**Example.**



**Remark 2.3.**  $l \neq u \neq k$ .

**Definition 2.4.** Define a set  $K(v, G) \subseteq V(G)$  as a set of vertices  $z \in V(G)$  such the path between  $z$  and  $l$  uses the edge  $\{w, u\}$ .

**Remark 2.5.**  $K(v, G)$  is the connected component of  $(V(G), E(G) \setminus \{w, u\})$ ,  $G$  without edge  $\{w, u\}$ , which contains  $v$ .

- Lemma 2.6.**
- (i) Paths between vertices from  $K(v, G)$  don't use the edge  $\{w, u\}$ .
  - (ii) Paths between vertices from  $V(G) \setminus K(v, G)$  don't use the edge  $\{w, u\}$ .
  - (iii) Paths joining a vertex from  $V(G) \setminus K(v, G)$  to a vertex from  $K(v, G)$  use the edge  $\{w, u\}$ .

**Proof.** Because  $G$  is a tree, there is a unique path between each pair of vertices, then it is obvious by remark 2.5. □

**Definition 2.7.** Define graphs

$$\begin{aligned} \bar{G} &= (V(G), E(G) \setminus \{\{w, u\}\} \cup \{\{w, l\}\}), \\ \tilde{G} &= (V(G), E(G) \setminus \{\{w, u\}\} \cup \{\{w, k\}\}). \end{aligned}$$

**Lemma 2.8.**  $\bar{G}$  and  $\tilde{G}$  are trees.

**Proof.** At first we show connectivity, let  $a, b \in V(G)$ , connect them with a path. If both are in  $K(v, G)$  or in  $V(G) \setminus K(v, G)$ , then by Lemma 2.6, the path in  $G$  uses only edges which are also in  $\bar{G}, \tilde{G}$ . Hence it is path also there.

Let  $a \in K(v, G)$  and  $b \in V(G) \setminus K(v, G)$ . We can see  $w \in K(v, G)$ , by Lemma 2.6 a path between  $a$  and  $w$ ,  $a = a_1, a_2, \dots, a_p = w$ , doesn't use  $\{w, u\}$  and all vertices of this path are in  $K(v, G)$ . If not, there is a path between vertices from  $K(v, G)$  and  $V(G) \setminus K(v, G)$  which doesn't use  $\{w, u\}$ , that is a contradiction with Lemma 2.6. Connect  $l$  and  $b$  with a path,  $l = b_1, b_2, \dots, b_q = b$ . It doesn't use  $\{w, u\}$  and all vertices are in  $V(G) \setminus K(v, G)$ . Then  $a = a_1, a_2, \dots, a_p = w$ ,  $l = b_1, b_2, \dots, b_q = b$  is a path between  $a, b$  in the graph  $\bar{G}$ , analogously for  $\tilde{G}$ .

Now we show that they don't contain a cycle, for contradiction suppose that  $\bar{G}$  contains a cycle  $C \subseteq \bar{G}$ . If  $C$  doesn't use the edge  $\{w, l\}$ , then  $C \subseteq G$ , but  $G$  is a tree, this is a contradiction. If  $C$  uses  $\{w, l\}$ , then there exists a path in  $G$  between  $w, l$ , which doesn't use the edge  $\{w, l\}$ . Then there exists a path in  $G$  between  $w, l$ , which doesn't use the edge  $\{w, u\}$ , but  $w \in K(v, G)$  and  $l \in V(G) \setminus K(v, G)$ , that is contradiction with Lemma 2.6. Analogously for  $\tilde{G}$ .  $\square$

We want to show that

$$s_H(G, f) \leq s_H(\bar{G}, f)$$

or

$$s_H(G, f) \leq s_H(\tilde{G}, f).$$

**Lemma 2.9.**

$$\begin{aligned} a, b \in K(v, G), & \quad \text{then } \rho_G(a, b) = \rho_{\bar{G}}(a, b) = \rho_{\tilde{G}}(a, b), \\ a, b \in V(G) \setminus K(v, G), & \quad \text{then } \rho_G(a, b) = \rho_{\bar{G}}(a, b) = \rho_{\tilde{G}}(a, b). \end{aligned}$$

**Proof.** A path in  $G$  between  $a, b$ , by Lemma 2.6, doesn't use  $\{u, w\}$ , hence it is a path in  $\bar{G}$  and  $\tilde{G}$  too, then the distance of  $a, b$  is the same in  $G, \bar{G}$  and  $\tilde{G}$ .  $\square$

**Definition 2.10.** Define subsets

$$F^+, F^-, F^0 \subseteq K(v, G) \times (V(G) \setminus K(v, G))$$

such that  $(a, b) \in F^+$  if a path between  $a, b$  uses the edge  $\{u, u^+\}$ .  $(a, b) \in F^-$  if a path between  $a, b$  uses the edge  $\{u, u^-\}$  and  $(a, b) \in F^0$  if a path between  $a, b$  doesn't use neither  $\{u, u^-\}$  nor  $\{u, u^+\}$ .

**Lemma 2.11.**  $F^+, F^-, F^0$  are pairwise disjoint and

$$F^+ \cup F^- \cup F^0 = K(v, G) \times (V(G) \setminus K(v, G)).$$

**Proof.** From the definition of  $F^+, F^-, F^0$  we have  $F^-$  and  $F^0, F^+$  and  $F^0$  are disjoint. Let  $(a, b) \in F^+ \cap F^-$ , then the path between  $a, b$  uses edges  $\{u, u^-\}, \{u, u^+\}$  and by lemma 2.6, it also uses the edge  $\{w, u\}$ . Hence it is a path which has a vertex of degree 3 and that is contradiction.  $\square$

**Lemma 2.12.** *Let  $x, \bar{x} \in K(v, G)$  and  $y, \bar{y} \in V(G) \setminus K(v, G)$  such that  $(x, y) \in F^+$  and  $(\bar{x}, \bar{y}) \in F^-$ . Then*

$$\begin{aligned} \rho_{\bar{G}}(x, y) + \rho_{\bar{G}}(\bar{x}, \bar{y}) &\geq \rho_G(x, y) + \rho_G(\bar{x}, \bar{y}), \\ \rho_{\bar{G}}(x, y) + \rho_{\bar{G}}(\bar{x}, \bar{y}) &\geq \rho_G(x, y) + \rho_G(\bar{x}, \bar{y}). \end{aligned}$$

*Moreover, both sides are equal, in the first inequality, if and only if  $y = l$  and, in the second inequality, if and only if  $\bar{y} = k$ .*

**Proof.** Let  $z$  denote the first common vertex of paths  $Q: l = y_1, y_2, \dots, y_s = k$  and  $P: y = x_1, x_2, \dots, x_m = x$ . Consider

$$i_m = \min\{i \mid \exists j \in \{1, \dots, m\}, y_i = x_j\}$$

and therefore  $z = y_{i_m}$ , let  $T$  be the path from  $z$  to  $l$ , we will show that  $z$  is the only one common vertex of  $T$  and  $P$ , vertices from  $P$  split into the 4 subpaths,  $P_1$  from  $y$  to  $z$ ,  $P_2$  from  $z$  to  $u$ , edge  $\{u, w\}$  and  $P_3$  from  $w$  to  $x$ . Vertices from  $P_1$  are not in  $Q$  (except for  $z$ ) from the definition of  $z$ . Vertices from  $P_2$  are not in  $T$  (except for  $z$ ) from the uniqueness of paths in trees and vertices from  $P_3$  belong to  $K(v, G)$  and every vertex of  $T$  belongs to  $V(G) \setminus K(v, G)$ . By composition of paths  $P_1, T, \{l, w\}, P_3$ , we get a path from  $y$  to  $x$  in the graph  $\bar{G}$ .

Let  $\bar{P}$  denote the path from  $\bar{y}$  to  $\bar{x}$ , analogously define  $\bar{z}$  as the first common vertex of paths  $\bar{P}$  and  $Q$  (first in the direction from  $\bar{y}$  to  $\bar{x}$ ). We split  $\bar{P}$  into the subpaths  $\bar{P}_1$  from  $\bar{y}$  to  $\bar{z}$ ,  $\bar{P}_2$  from  $\bar{z}$  to  $u$ , edge  $\{u, w\}$  and  $\bar{P}_3$  from  $u$  to  $\bar{x}$ . Let  $\bar{T}$  be the path from  $u$  to  $l$ , analogously we get that  $u$  is the only one common vertex of  $\bar{P}$  and  $\bar{T}$ . Hence  $\bar{P}_1, \bar{P}_2, \bar{T}, \{l, w\}, \bar{P}_3$  is a path between  $\bar{y}, \bar{x}$  in the graph  $\bar{G}$ .

And for paths from  $u$  to  $z$  and from  $u$  to  $\bar{z}$ ,  $u$  is the only one common vertex, by uniqueness of path in trees.

Now we can calculate

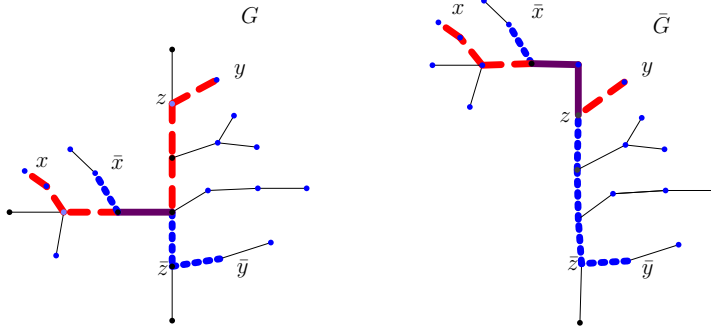
$$\begin{aligned} \rho_G(x, y) &= \rho_G(x, w) + 1 + \rho_G(u, z) + \rho_G(z, y), \\ \rho_G(\bar{x}, \bar{y}) &= \rho_G(\bar{x}, w) + 1 + \rho_G(u, \bar{z}) + \rho_G(\bar{z}, \bar{y}), \\ \rho_{\bar{G}}(x, y) &= \rho_G(x, w) + 1 + \rho_G(l, z) + \rho_G(z, y), \\ \rho_{\bar{G}}(\bar{x}, \bar{y}) &= \rho_G(\bar{x}, w) + 1 + \rho_G(l, z) + \rho_G(z, u) + \rho_G(u, \bar{z}) + \rho_G(\bar{z}, \bar{y}), \end{aligned}$$

hence

$$\rho_{\bar{G}}(\bar{x}, \bar{y}) + \rho_{\bar{G}}(x, y) = \rho_G(\bar{x}, \bar{y}) + \rho_G(x, y) + 2\rho_G(l, z).$$

Now we get our inequality and we see that both are equal if and only if  $l = z$ . But  $l$  is a leaf, hence  $z$  is a leaf, then  $y = z = l$ . For  $\bar{G}$  analogously.  $\square$

**Example.** Paths between  $x, y$  and  $\bar{x}, \bar{y}$  in graphs  $G$  and  $\bar{G}$ .



**Lemma 2.13.** *Let  $(x, y) \in F^0$  then*

$$\begin{aligned} \rho_{\bar{G}}(x, y) &> \rho_G(x, y), \\ \rho_{\tilde{G}}(x, y) &> \rho_G(x, y). \end{aligned}$$

**Proof.** Let  $P$  be a path from  $x$  to  $y$  and  $Q$  be a path from  $l$  to  $k$  in  $G$ , for  $P$  and  $Q$ ,  $u$  is the only one common vertex because  $(x, y) \in F^0$ . Hence  $x \rightarrow w - l \rightarrow u \rightarrow y$  is a path in  $\bar{G}$ , where paths of type  $a \rightarrow b$  are subpaths of  $P$  and  $Q$  and  $-$  denotes an edge. Now we can calculate the following

$$\rho_{\bar{G}}(x, y) = \rho_G(x, u) + 1 + \rho_G(l, u) + \rho_G(u, y) = \rho_G(x, y) + \rho_G(l, u)$$

and from  $l \neq u$  we have inequality.

For  $\tilde{G}$  analogously. □

**Lemma 2.14.**

$$\begin{aligned} \rho_{\bar{G}}(x, y) &> \rho_G(x, y) \quad \text{for } (x, y) \in F^-, \\ \rho_{\tilde{G}}(x, y) &> \rho_G(x, y) \quad \text{for } (x, y) \in F^+. \end{aligned}$$

**Proof.** We will prove the first inequality. As well as in lemma 2.12 denote  $z$  the first common vertex of paths from  $y$  to  $x$  and from  $k$  to  $l$ , formally we can define it as well as in lemma 2.12. Now we consider a path  $x \rightarrow w - l \rightarrow u \rightarrow z \rightarrow y$ . Hence

$$\begin{aligned} \rho_{\bar{G}}(x, y) &= \rho_G(x, w) + 1 + \rho_G(l, u) + \rho_G(u, z) + \rho_G(z, y) \\ &= \rho_G(x, y) + \rho_G(l, u) \end{aligned}$$

and from  $l \neq u$  we have inequality.

For second inequality analogously. □

**Definition 2.15.** Let  $G$  be a tree and  $H$  be a graph such that

$$|V(G)| = |V(H)|$$

and

$$f: V(H) \rightarrow V(G)$$

is a pseudoordering, we define a set

$$L = \{(x, y) \in K(v, G) \times (V(G) \setminus K(v, G)) \mid x \sim_{H,f} y\},$$

where  $K(v, G)$  is the set from Definition 2.4.

**Lemma 2.16.** *Let  $G$  be a tree and  $H$  be a graph such that,  $|V(G)| = |V(H)|$  and*

$$f: V(H) \rightarrow V(G)$$

*is a pseudoordering. Then*

$$s_H(f, \bar{G}) \geq s_H(f, G)$$

*or*

$$s_H(f, \tilde{G}) \geq s_H(f, G),$$

*the first case occurs when*

$$|L \cap F^+| \leq |L \cap F^-|,$$

*the second case occurs when*

$$|L \cap F^+| \geq |L \cap F^-|.$$

**Proof.** Denote  $n^+ = |L \cap F^+|$ ,  $n^- = |L \cap F^-|$ ,  $m = |L \cap F^0|$ ,

$$\bar{m} = \frac{|\{(x, y) \in (K(v, G)^2) \cup ((V(G) \setminus K(v, G))^2) \mid x \sim_{H,f} y\}|}{2},$$

where square  $K(v, G)^2$  means  $K(v, G) \times K(v, G)$ .  $\bar{m}$  is number of edges  $\{x, y\} \in E(H)$ , which satisfy that  $f(x)$  and  $f(y)$  lie in the same component of

$$(V(G), E(G) \setminus \{w, u\}).$$

Let  $n^+ \geq n^-$ , the second case is analogous, we rearrange the sum  $s_H(f, G)$  in this way

$$\begin{aligned} s_H(f, G) &= \sum_{i=1}^{n^-} (\rho_G(x_i, y_i) + \rho_G(\bar{x}_i, \bar{y}_i)) + \sum_{i=n^-+1}^{n^+} \rho_G(x_i, y_i) \\ &\quad + \sum_{i=1}^m \rho_G(a_i, b_i) + \sum_{i=1}^{\bar{m}} \rho_G(c_i, d_i), \end{aligned}$$

where

$$(x_i, y_i) \in F^+, \quad (\bar{x}_i, \bar{y}_i) \in F^-, \quad (a_i, b_i) \in F^0,$$

$$(c_i, d_i) \in \{(x, y) \in (K(v, G)^2) \cup ((V(G) \setminus K(v, G))^2) \mid x \sim_{H,f} y\}.$$

Now, by Lemma 2.12

$$\rho_G(x_i, y_i) + \rho_G(\bar{x}_i, \bar{y}_i) \leq \rho_{\bar{G}}(x_i, y_i) + \rho_{\bar{G}}(\bar{x}_i, \bar{y}_i),$$

by Lemma 2.14

$$\rho_G(x_i, y_i) \leq \rho_{\bar{G}}(x_i, y_i),$$



by Lemma 2.13

$$\rho_G(a_i, b_i) \leq \rho_{\bar{G}}(a_i, b_i)$$

and by Lemma 2.9

$$\rho_G(c_i, d_i) = \rho_{\bar{G}}(c_i, d_i).$$

Hence

$$\begin{aligned} s_H(f, G) &\leq \sum_{i=1}^{n^-} (\rho_{\bar{G}}(x_i, y_i) + \rho_{\bar{G}}(\bar{x}_i, \bar{y}_i)) \\ &\quad + \sum_{i=n^-+1}^{n^+} \rho_{\bar{G}}(x_i, y_i) + \sum_{i=1}^m \rho_{\bar{G}}(a_i, b_i) + \sum_{i=1}^{\bar{m}} \rho_{\bar{G}}(c_i, d_i) \\ &= s_H(f, \tilde{G}). \end{aligned}$$

□

**Lemma 2.17.** *Let  $G$  be a tree and  $H$  be a graph such that,  $|V(G)| = |V(H)|$  and  $f: V(H) \rightarrow V(G)$  is a pseudoordering. Then there exists a pseudoordering*

$$g: V(H) \rightarrow \{x_1, x_2, \dots, x_{|V(G)|}\} = V(P_{|V(G)|-1}) \quad \text{such that}$$

$$s_H(f, G) \leq s_H(g, P_{|V(G)|-1}).$$

**Proof.** We denote

$$\alpha(G) = \sum_{\substack{v \in V(G) \\ \deg_G v \geq 3}} \deg_G v,$$

from the definition of  $u, l$  and  $k$  we know that  $\deg_G u \geq 3$  and  $\deg_G l = \deg_G k = 1$ . From the construction of  $\bar{G}$  and  $\tilde{G}$  we have  $\deg_{\bar{G}} u = \deg_{\tilde{G}} u \leq \deg_G u$ ,  $\deg_{\bar{G}} l = \deg_{\tilde{G}} l = \deg_G l = 1$  and all other vertices have the same degree as before. Hence

$$\alpha(\bar{G}) < \alpha(G),$$

$$\alpha(\tilde{G}) < \alpha(G).$$

Let  $S$  be a tree, which is not a path, we choose any three pairwise distinct leaves in  $V(S)$  and define  $S^*$  as one of graphs  $\bar{S}, \tilde{S}$ , which satisfy  $s_H(f, S^*) \geq s_H(f, S)$ . Denote  $G_0 = G$  and for  $i \geq 0$  denote  $G_{i+1} = G_i^*$  if  $G_i$  is not a path, otherwise define  $G_{i+1} = G_i$ . For contradiction we assume that the tree  $G_i$  is not a path for every  $i \in \mathbb{N}_0$ . We know  $\alpha(G_i) \in \mathbb{N}_0$  for every  $i$  and

$$\alpha(G_{i+1}) \leq \alpha(G_i) - 1,$$

hence

$$\alpha(G_{\alpha(G_0)+1}) \leq \alpha(G_0) - \alpha(G_0) - 1 = -1$$

and this is contradiction. Therefore there exists some  $j$  such that  $G_j$  is a path, from Lemma 2.16 we get

$$s_H(f, G_{i+1}) \geq s_H(f, G_i)$$

and hence

$$s_H(f, G_j) \geq s_H(f, G).$$

□

**Theorem 2.18.** *Let  $G$  and  $H$  be graphs such that  $G$  is connected,  $|V(G)| = |V(H)|$  and  $f: V(H) \rightarrow V(G)$  is a pseudoordering, then there exists a pseudordering*

$$g: V(H) \rightarrow \{x_1, x_2, \dots, x_{|V(G)|}\} = V(P_{|V(G)|-1}) \quad \text{such that}$$

$$s_H(f, G) \leq s_H(g, P_{|V(G)|-1}).$$

**Proof.** Let  $K$  be any spanning tree of  $G$ ,  $x, y \in V(G)$ , we connect  $x$  and  $y$  with a path in graph  $K$ , this path is also a path in  $G$ . Hence

$$\rho_G(x, y) \leq \rho_K(x, y)$$

for every  $x, y$ , hence

$$s_H(f, G) \leq s_H(f, K),$$

by Lemma 2.17 there exists a pseudoordering

$$g: V(H) \rightarrow \{x_1, x_2, \dots, x_{|V(G)|}\} = V(P_{|V(G)|-1}) \quad \text{such that}$$

$$s_H(f, G) \leq s_H(f, K) \leq s_H(g, P_{|V(G)|-1}).$$

□

**Corollary 2.19.** *Let  $G$  and  $H$  be graphs such that  $G$  is connected,  $|V(G)| = |V(H)|$ , then*

$$h_H^+(G) \leq h_H^+(P_{|V(G)|-1}).$$

### 3. GRAPHS WITH A MAXIMAL UPPER H-HAMILTONIAN NUMBER

In this section we will prove that if in Corollary 2.19 the graph  $H$  is connected, then in the inequality in Corollary 2.19 both sides are equal.

**Remark 3.1.** For easier writing, we will denote vertices of  $H$  the same as vertices of  $G$ , we will rename them in this way  $v \in H \mapsto f(v)$ . We can naturally see it as graph with two sets of edges.

In inequalities in Lemma 2.16 both sides are equal under specific conditions, if  $L \cap F^0 \neq \emptyset$ , then in Lemma 2.13 there is a strict inequality and then also the same happens in Theorem 2.18.

If  $(L \setminus K(v, G) \times \{l\}) \cap F^+ \neq \emptyset$ , then in Lemma 2.12 there is a strict inequality and then also the same happens in Theorem 2.18. Analogously if

$$(L \setminus K(v, G) \times \{k\}) \cap F^- \neq \emptyset.$$

Overall we get that the only nontrivial case is

$$(1) \quad L \subseteq K(v, G) \times \{k, l\}.$$

**Remark 3.2.** Remark 3.1 holds for every triple of distinct leaves  $k, l, v$  in  $G$ .

**Lemma 3.3.** *Let  $G$  be a tree,  $H$  connected graph such that  $|V(G)| = |V(H)|$  and  $f: V(H) \rightarrow V(G)$  is a pseudoordering, which satisfy*

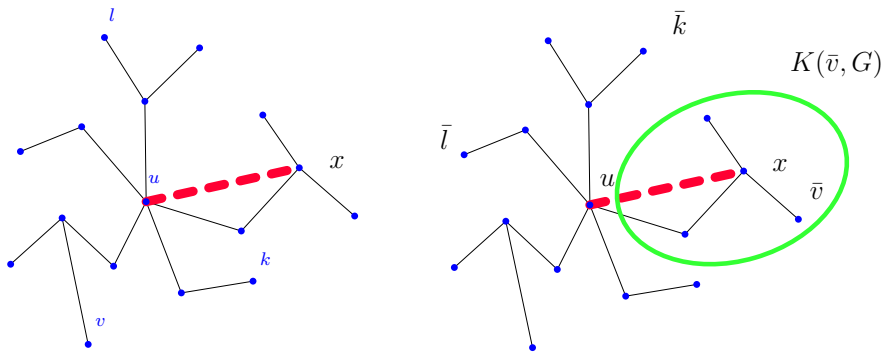
$$s_H(f, G) = h_H^+(P_{|V(G)|-1}),$$

*then  $G$  is path.*

**Proof.** For contradiction suppose that  $G$  is not a path, then there exist three pairwise distinct leaves  $k, l, v$ , we denote in the same way as before, vertex  $u$  and set of vertices  $K(v, G)$ . Because graph  $H$  is connected there exists a vertex  $x$  such that  $\{u, x\} \in E(H)$ . Let  $X \subseteq V(G)$  be a set of vertices of components of graph  $G \setminus u$ , containing  $x$ .  $G \setminus u$  has, by definition of  $u$ , at least 3 components. Let now  $\bar{v}$  be an arbitrary leaf (leaf in  $G$ ) in  $X$ . Choose  $\bar{k}, \bar{l}$  as arbitrary leaves in pairwise distinct components of  $G \setminus u$  and different from  $X$ .

Now  $(x, u) \in \bar{L}$ , where  $\bar{L}$  is alternative of  $L$  for  $\bar{k}, \bar{l}, \bar{v}$  and by Remark 3.1 for  $\bar{k}, \bar{l}, \bar{v}$  and by  $k \neq u \neq l$  we get contradiction. □

**Example.** We show the idea of the last proof in the following picture.



**Remark 3.4.** Let  $G$  be a graph with a maximal  $H$ -Hamiltonian number, then every spanning tree of  $G$  has a maximal  $H$ -Hamiltonian number, therefore every spanning tree is a path. We will show that the only graphs with this property are cycles and paths.

**Lemma 3.5.** *Let  $G$  be a connected graph such that  $|V(G)| \geq 2$ , then there is a vertex, which is not an articulation point.*

**Proof.** Consider a block-cut tree of  $G$  and a block  $B$ , which is a leaf of the block-cut tree or if this tree has only one vertex, then  $B = G$ .  $B$  is, by definition of a block, 2-connected. Because  $B$  is leaf we get that in  $B$  there is only one articulation and in  $B$  there are at least 2 vertices. Hence in  $B$  there is at least one vertex, which is not an articulation point. □

**Lemma 3.6.** *Let  $G$  be a finite connected graph such that  $|V(G)| \geq 2$  and every spanning tree of  $G$  is a path, then  $G$  is a path or a cycle.*

**Proof.** We will prove it by induction with respect to the number of vertices. Let  $n$  be the number of vertices, for  $n = 2$  and  $n = 3$  it is obviously true. Let it be true for  $n \geq 3$ , let  $G$  be a graph with  $n + 1$  vertices such that every spanning tree of  $G$  is a path. Let  $v \in V(G)$  be a vertex, which is not an articulation point, by lemma 3.5 it exists. We denote  $G'$  the subgraph induced by the set of vertices  $V(G) \setminus \{v\}$ .  $G'$  is connected, we will show that every spanning tree of  $G'$  is a path. Let there exist a spanning tree which is not a path, let  $u \in V(G)$  be a vertex such that  $\{v, u\} \in E(G)$ . Now when we add this edge to the spanning tree, we get a spanning tree of  $G$ , which is not a path and it is a contradiction. By induction hypothesis  $G'$  is a path or a cycle, we denote  $A = \{u \in V(G) | \{v, u\} \in E(G)\}$ . For contradiction we assume  $G'$  is a cycle and let  $u \in A$ , in  $G'$  be an edge  $e$  such that  $u$  is not incident to  $e$ . Consider the subgraph  $B$  of  $G$ ,  $B = (V(G), E(G') \setminus e \cup \{v, u\})$ , and this is a spanning tree of  $G$  which is not a path, contradiction.

Therefore  $G'$  is a path, let  $x, y$  be endpoints of this path, for contradiction we assume that there exists some another vertex  $u \in A$ . Hence  $G'$  together with  $\{u, v\}$  form a spanning tree which is not a path. Hence  $A \subseteq \{x, y\}$ , because  $G$  is connected we get also  $A \neq \emptyset$ . Finally there are the two cases for  $G$ , if  $|A| = 1$ , then  $G$  will be a path and if  $|A| = 2$ , then  $G$  will be a cycle.  $\square$

**Theorem 3.7.** *Let  $G$  and  $H$  be connected finite graphs such that  $|V(G)| = |V(H)|$ , then*

$$h_H^+(G) \leq h_H^+(P_{|V(G)|-1}),$$

moreover, both sides are equal if and only if  $G$  is a path.

**Proof.** The first part follows from Theorem 2.18, let  $G$  be a graph,  $f$  be a pseudoordering such that

$$s_H(f, G) = h_H^+(G) = h_H^+(P_{|V(G)|-1}).$$

From the proof of Theorem 2.18 we know that every spanning tree also satisfies the equation above. Hence, by Lemma 3.3, every spanning tree of  $G$  is a path. By Lemma 3.6  $G$  is a path or a cycle, for contradiction we assume, that it is a cycle. We denote  $n = |V(G)|$ , we will show that there are two vertices  $v, u \in V(G)$  such that  $v \sim_{H,f} u$  and  $\rho_G(u, v) < \frac{n}{2}$ .

Because  $G$  is cycle,  $|V(H)| = n \geq 3$  and  $H$  is connected we see that there is a vertex of degree at least 2. Let  $v$  be a vertex such that  $\deg_H(v) \geq 2$ , there exists at least two vertices  $u$  such that  $v \sim_{H,f} u$ . There exists at most one vertex such that  $\rho_G(u, v) \geq \frac{n}{2}$ , hence at least one of them satisfies  $\rho_G(u, v) < \frac{n}{2}$ .

Now we connect  $v$  and  $u$  with a shorter path in  $G$ . Let  $e$  be some edge on this path, we define a graph  $\bar{G} = (V(G), E(G) \setminus e)$ , it is a path, where every distance is greater or equal as in  $G$ . But  $\rho_G(u, v) < \rho_{\bar{G}}(u, v)$  and then

$$s_H(f, \bar{G}) = s_H(f, \bar{G}) > h_H^+(P_{|V(G)|-1}),$$

and this is contradiction with Theorem 2.18.  $\square$

## 4. CONCLUSION

When we use following equations which can be found for example in [2, Theorem 1.3] and [2, Corollary 2.2]

$$h^+(P_{|V(G)|-1}) = \left\lfloor \frac{|V(G)|^2}{2} \right\rfloor, \quad t^+(P_{|V(G)|-1}) = \left\lfloor \frac{|V(G)|^2}{2} \right\rfloor - 1.$$

This result is also calculated in [1] and when we use Theorem 3.7 for  $H = P_{|V(G)|-1}$  and for  $H = C_{|V(G)|}$  we get the following theorem.

**Theorem 4.1** ([2]).

$$h^+(G) \leq \left\lfloor \frac{|V(G)|^2}{2} \right\rfloor, \quad t^+(G) \leq \left\lfloor \frac{|V(G)|^2}{2} \right\rfloor - 1.$$

Moreover, both sides are equal if and only if  $G$  is a path.

First part is [2, Corollary 2.2] and second part is [2, Theorem 4.2]. Now we can see, that Theorem 3.7 is generalization of Theorem 4.1 which is from article [2].

## REFERENCES

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DEPARTMENT OF MATHEMATICS AND STATISTICS,  
 FACULTY OF SCIENCE, MASARYK UNIVERSITY,  
 KOTLÁŘSKÁ 2, 611 37 BRNO, CZECH REPUBLIC  
*E-mail*: 451859@mail.muni.cz