

CYCLIC AND λ -CONSTACYCLIC CODES OVER THE RING

$$\mathbb{Z}_5[u, v] / \langle u^2 - u, v^2, uv, vu \rangle$$

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ABSTRACT. In this study, unitary elements and related elements are determined on two variable rings with coefficients of \mathbb{Z}_5 . For the $u^2 = u, v^2 = 0$ and $u \cdot v = v \cdot u = 0$ states, λ -constacyclic codes and the types of codes with their gray images were determined over $\lambda = (1 + 3u), (2 + 4u)$ and 4 unitary elements on the $\mathbb{Z}_5[u, v] / \langle u^2 - u, v^2, uv, vu \rangle$ ring. It has been shown that codes with $[5n, k, d_H]$ parameter are obtained on the \mathbb{Z}_5 object.

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1. INTRODUCTION

Cyclic codes, constacyclic codes, quasi cyclic codes, negacyclic codes and skew cyclic codes were studied in one-variable and two-variable rings with coefficients of \mathbb{Z}_{p^k} field being p a prime number and k an integer. Most of these studies have been codes in the literature corresponding to \mathbb{Z}_p prime fields for $k = 1$. Basic information in coding theory, parameters of codes and code definitions are given in the book of Steven Roman [1], which is a general reference. In one-variable rings; Constacyclic codes for unitary element $1 + u$ in ring $\mathbb{F}_2 + u\mathbb{F}_2$ with 4 elements whose coefficients are in binary field Qian J. and his team in [2]. Study in [2], it has been shown that $(1 + u)$ -constacyclic codes correspond to cyclic codes on the field, thanks to the Gray transform between the four elements ring and the binary body. In [3], with a different method, the work in this four elements ring was transferred to the eight elements $\mathbb{F}_2 + u\mathbb{F}_2 + u^2\mathbb{F}_2$ ring and similar results were obtained. Previous studies have addressed constacyclic strains in bivariate rings, similar to studies in univariate rings. The study titled "On some special codes over $\mathbb{F}_3 + v\mathbb{F}_3 + u\mathbb{F}_3 + u^2\mathbb{F}_3$ " written by M.Özkan, which we frequently use in this article, and the constacyclic codes on the rings with the coefficients in the ternary field and their images in the \mathbb{F}_3 field are presented in [4]. In [5], a class of constacyclic codes in the bivariate ring with

coefficients of \mathbb{Z}_4 for $p = 2$ and $k = 2$ cases is given by H. Islam. In another article, Gray images of constacyclic codes for ring $\mathbb{F}_2 + u_1\mathbb{F}_2 + u_2\mathbb{F}_2$ for bivariate variables u_1 and u_2 and which codes they are have been studied in [6]. In [7], the images of the codes under Gray transform on bivariate rings with \mathbb{Z}_3 coefficient published by Timothy Kom and his team are given. In this study, a new Gray transform is defined using the ring presented in [7]. A different perspective has been gained for the codes under the Gray transformation and new codes have been written.

2. PRELIMINARIES

Let $S = \mathbb{Z}_5[u, v]/\langle u^2 - u, v^2, uv, vu \rangle$ and $\mathbb{Z}_5 = \{\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}\}$. Then $S = \{a + ub + vc : a, b, c \in \mathbb{Z}_5\}$ is a commutative ring with cardinality 125 and characteristic 5. The set of units of the ring is $vI = \{1, 4, 1 + 3u, 2 + 4u\} = \{\lambda \in I : \lambda^2 = 1\}$. In this study, 4 of the unitary elements of the ring were examined. The ring S contains more than one maximal ideal. Hence, it is a finite non-chain ring. Also, $S = \mathbb{Z}_5[u, v]/\langle u^2 - u, v^2, uv, vu \rangle \cong \mathbb{Z}_5 + u\mathbb{Z}_5 + v\mathbb{Z}_5$ with $u^2 = u, v^2 = 0$ and $u \cdot v = v \cdot u = 0$. The definitions to be used in this study are given below.

Definition 1. A linear code C over S of length n is a S -submodule of S^n . An element of C is named a codeword. A cyclic code C of length n over S is a linear code with the characteristic that if $c = (c_0, c_1, c_2, c_3, \dots, c_{n-1}) \in C$, then $\sigma(c) = (c_{n-1}, c_0, c_1, c_2, \dots, c_{n-2}) \in C$. σ is named cyclic shift operator from S^n to S^n .

Definition 2. A linear code C of length n over S is λ -constacyclic code if $c = (c_0, c_1, c_2, c_3, \dots, c_{n-1}) \in C$, then $\gamma_{(\lambda)}(c) = (\lambda c_{n-1}, c_0, c_1, c_2, c_3, \dots, c_{n-2}) \in C$, where λ is a unit in S . $\gamma_{(\lambda)}$ is called λ -constacyclic shift operator from S^n to S^n .

Definition 3. Let $a \in \mathbb{Z}_5^{3n}$ with $a = (a_0, a_1, a_2, \dots, a_n, \dots, a_{2n}, \dots, a_{3n-1}) = (a^{(0)}|a^{(1)}|a^{(2)})$, where $a^{(i)} \in \mathbb{Z}_5^n$ for $i = 0, 1, 2$ and $|$ is the usual vector concatenation. Let ρ be a map from \mathbb{Z}_5^{3n} to \mathbb{Z}_5^{3n} defined by $\rho(a) = (\sigma(a^{(0)})|\sigma(a^{(1)})|\sigma(a^{(2)}))$ where σ is a cyclic shift operator from \mathbb{Z}_5^n to \mathbb{Z}_5^n . A code C of length $3n$ over \mathbb{Z}_5 is called a quasi-cyclic code of index 3 if $\rho(C) = C$.

Proposition 1. A subset C of S^n is a $[n, d]$ -cyclic code if and only if its polynomial representation is an ideal of $S_n = S[x]/\langle x^n - 1 \rangle$.

Proposition 2. A subset C of S^n is a constacyclic code of length n if and only if its polynomial representation is an ideal of $S_{n, \lambda} = S[x]/\langle x^n - \lambda \rangle$.

3. GRAY MAP AND CYCLIC CODES OVER S

In this section, we introduce a Gray map Γ on the ring S and consider the algebraic structures of cyclic codes over the ring S .

In order to connect the structure of the ring S with \mathbb{Z}_5^3 . We define the Gray map Γ ;

$$\Gamma : S \rightarrow \mathbb{Z}_5^3$$

$$a + ub + vc \rightarrow \Gamma(a + ub + vc) = (a + 4b, b, c)$$

where $a + ub + vc \in S$ and $a, b, c \in \mathbb{Z}_5$. From the definition, we observe that

$$\begin{aligned} \Gamma(0) &= (0, 0, 0), \Gamma(1) = (1, 0, 0), \Gamma(2) = (2, 0, 0), \Gamma(3) = (3, 0, 0), \Gamma(4) = (4, 0, 0), \\ \Gamma(u) &= (4, 1, 0), \Gamma(2u) = (3, 2, 0), \Gamma(3u) = (2, 3, 0), \Gamma(4u) = (1, 4, 0), \\ \Gamma(v) &= (0, 0, 1), \Gamma(2v) = (0, 0, 2), \Gamma(3v) = (0, 0, 3), \Gamma(4v) = (0, 0, 4), \\ \Gamma(1+u) &= (0, 1, 0), \Gamma(1+2u) = (4, 2, 0), \Gamma(1+3u) = (3, 3, 0), \Gamma(1+4u) = (2, 4, 0), \\ \Gamma(2+u) &= (1, 1, 0), \Gamma(2+2u) = (0, 2, 0), \Gamma(2+3u) = (4, 3, 0), \Gamma(2+4u) = (3, 4, 0), \\ \Gamma(3+u) &= (2, 1, 0), \Gamma(3+2u) = (1, 2, 0), \Gamma(3+3u) = (0, 3, 0), \Gamma(3+4u) = (4, 4, 0), \\ \Gamma(4+u) &= (3, 1, 0), \Gamma(4+2u) = (2, 2, 0), \Gamma(4+3u) = (1, 3, 0), \Gamma(4+4u) = (0, 4, 0), \\ \Gamma(1+v) &= (1, 0, 1), \Gamma(1+2v) = (1, 0, 2), \Gamma(1+3v) = (1, 0, 3), \Gamma(1+4v) = (1, 0, 4), \\ \Gamma(2+v) &= (2, 0, 1), \Gamma(2+2v) = (2, 0, 2), \Gamma(2+3v) = (2, 0, 3), \Gamma(2+4v) = (2, 0, 4), \\ \Gamma(3+v) &= (3, 0, 1), \Gamma(3+2v) = (3, 0, 2), \Gamma(3+3v) = (3, 0, 3), \Gamma(3+4v) = (3, 0, 4), \\ \Gamma(4+v) &= (4, 0, 1), \Gamma(4+2v) = (4, 0, 2), \Gamma(4+3v) = (4, 0, 3), \Gamma(4+4v) = (4, 0, 4), \\ \Gamma(u+v) &= (4, 1, 1), \Gamma(u+2v) = (4, 1, 2), \Gamma(u+3v) = (4, 1, 3), \Gamma(u+4v) = (4, 1, 4), \\ \Gamma(2u+v) &= (3, 2, 1), \Gamma(2u+2v) = (3, 2, 2), \Gamma(2u+3v) = (3, 2, 3), \Gamma(2u+4v) = (3, 2, 4), \\ \Gamma(3u+v) &= (2, 3, 1), \Gamma(3u+2v) = (2, 3, 2), \Gamma(3u+3v) = (2, 3, 3), \Gamma(3u+4v) = (2, 3, 4), \\ \Gamma(4u+v) &= (1, 4, 1), \Gamma(4u+2v) = (1, 4, 2), \Gamma(4u+3v) = (1, 4, 3), \Gamma(4u+4v) = (1, 4, 4), \\ \Gamma(1+u+v) &= (0, 1, 1), \Gamma(1+u+2v) = (0, 1, 2), \Gamma(1+u+3v) = (0, 1, 3), \Gamma(1+u+4v) = (0, 1, 4), \\ \Gamma(1+2u+v) &= (4, 2, 1), \Gamma(1+2u+2v) = (4, 2, 2), \Gamma(1+2u+3v) = (4, 2, 3), \Gamma(1+2u+4v) = (4, 2, 4), \\ \Gamma(1+3u+v) &= (3, 3, 1), \Gamma(1+3u+2v) = (3, 3, 2), \Gamma(1+3u+3v) = (3, 3, 3), \Gamma(1+3u+4v) = (3, 3, 4), \\ \Gamma(1+4u+v) &= (2, 4, 1), \Gamma(1+4u+2v) = (2, 4, 2), \Gamma(1+4u+3v) = (2, 4, 3), \Gamma(1+4u+4v) = (2, 4, 4), \end{aligned}$$

$$\begin{aligned}
\Gamma(2+u+v) &= (1, 1, 1), \Gamma(2+u+2v) = (1, 1, 2), \Gamma(2+u+3v) = (1, 1, 3), \Gamma(2+u+4v) = (1, 1, 4), \\
\Gamma(2+2u+v) &= (0, 2, 1), \Gamma(2+2u+2v) = (0, 2, 2), \Gamma(2+2u+3v) = (0, 2, 3), \Gamma(2+2u+4v) = (0, 2, 4), \\
\Gamma(2+3u+v) &= (4, 3, 1), \Gamma(2+3u+2v) = (4, 3, 2), \Gamma(2+3u+3v) = (4, 3, 3), \Gamma(2+3u+4v) = (4, 3, 4), \\
\Gamma(2+4u+v) &= (3, 4, 1), \Gamma(2+4u+2v) = (3, 4, 2), \Gamma(2+4u+3v) = (3, 4, 3), \Gamma(2+4u+4v) = (3, 4, 4), \\
\Gamma(3+u+v) &= (2, 1, 1), \Gamma(3+u+2v) = (2, 1, 2), \Gamma(3+u+3v) = (2, 1, 3), \Gamma(3+u+4v) = (2, 1, 4), \\
\Gamma(3+2u+v) &= (1, 2, 1), \Gamma(3+2u+2v) = (1, 2, 2), \Gamma(3+2u+3v) = (1, 2, 3), \Gamma(3+2u+4v) = (1, 2, 4), \\
\Gamma(3+3u+v) &= (0, 3, 1), \Gamma(3+3u+2v) = (0, 3, 2), \Gamma(3+3u+3v) = (0, 3, 3), \Gamma(3+3u+4v) = (0, 3, 4), \\
\Gamma(3+4u+v) &= (4, 4, 1), \Gamma(3+4u+2v) = (4, 4, 2), \Gamma(3+4u+3v) = (4, 4, 3), \Gamma(3+4u+4v) = (4, 4, 4), \\
\Gamma(4+u+v) &= (3, 1, 1), \Gamma(4+u+2v) = (3, 1, 2), \Gamma(4+u+3v) = (3, 1, 3), \Gamma(4+u+4v) = (3, 1, 4), \\
\Gamma(4+2u+v) &= (2, 2, 1), \Gamma(4+2u+2v) = (2, 2, 2), \Gamma(4+2u+3v) = (2, 2, 3), \Gamma(4+2u+4v) = (2, 2, 4), \\
\Gamma(4+3u+v) &= (1, 3, 1), \Gamma(4+3u+2v) = (1, 3, 2), \Gamma(4+3u+3v) = (1, 3, 3), \Gamma(4+3u+4v) = (1, 3, 4), \\
\Gamma(4+4u+v) &= (0, 4, 1), \Gamma(4+4u+2v) = (0, 4, 2), \Gamma(4+4u+3v) = (0, 4, 3), \Gamma(4+4u+4v) = (0, 4, 4).
\end{aligned}$$

It can be easily checked that Γ is bijective. The map Γ can be extended in a natural way to S^n component-wise. For $q = (q_0, q_1, \dots, q_{n-1}) \in S^n$, Γ can be defined as follows:

$$\Gamma : S^n \rightarrow \mathbb{Z}_5^{3n}$$

$$\Gamma(q_0, q_1, \dots, q_{n-1}) = (a_0+4b_0, a_1+4b_1, \dots, a_{n-1}+4b_{n-1}, b_0, b_1, \dots, b_{n-1}, c_0, c_1, \dots, c_{n-1})$$

where $q_i = a_i + ub_i + vc_i \in S$ and $a_i, b_i, c_i \in \mathbb{Z}_5$ for $i = 0, 1, \dots, n-1$.

Let C be a linear code of length n over S . For any $r = (r_0, r_1, \dots, r_{n-1}) \in C$ the Hamming weight $w_H(C)$ of a code C is the smallest weight among all its non-zero codewords. For $r = (r_0, r_1, \dots, r_{n-1})$ and $r' = (r'_0, r'_1, \dots, r'_{n-1})$ in C , the Hamming distance between r and r' is defined by $d_H(r, r') = w_H(r - r')$ and the Hamming distance for a code C is defined by $d_H(C) = \min\{d_H(r, r') | r, r' \in C\}$.

The Lee weight of any element $r = (r_0, r_1, \dots, r_{n-1}) \in S^n$ is defined by $w_L(r) = \sum_{i=0}^{n-1} w_L(r_i)$, where $w_L(r_i) = w_H(a_i + 4b_i, b_i, c_i)$ for $r_i = a_i + ub_i + vc_i \in S, i = 0, 1, \dots, n-1$. The Lee distance for the code C is defined by

$d_L(C) = \min\{d_L(r, r') | r \neq r', \forall r, r' \in C\}$, where $d_L(r, r')$ is the Lee distance between r and r' defined by $d_L(r, r') = w_L(r - r')$.

Theorem 1. *The Gray map $\Gamma : S^n \rightarrow \mathbb{Z}_5^{3n}$ is a distance preserving \mathbb{Z}_5 -linear map from S^n (Lee distance, d_L) to \mathbb{Z}_5^{3n} (Hamming distance, d_H).*

Proof. Let $q = (q_0, q_1, \dots, q_{n-1}), k = (k_0, k_1, \dots, k_{n-1}) \in S^n$, where $q_i = a_i + ub_i + vc_i$, $k_i = e_i + uf_i + vg_i \in S$ for $i = 0, 1, \dots, n-1$ and $\beta \in \mathbb{Z}_5$. Then $\Gamma(q+k) = \Gamma(q_0+k_0, q_1+k_1, \dots, q_{n-1}+k_{n-1}) = (a_0+e_0+4(b_0+f_0), \dots, a_{n-1}+e_{n-1}+4(b_{n-1}+f_{n-1}), b_0+f_0, \dots, b_{n-1}+f_{n-1}, c_0+g_0, \dots, c_{n-1}+g_{n-1}) = (a_0+4b_0, \dots, a_{n-1}+4b_{n-1}, b_0, \dots, b_{n-1}, c_0, \dots, c_{n-1}) + (e_0+4f_0, \dots, e_{n-1}+4f_{n-1}, f_0, \dots, f_{n-1}, g_0, \dots, g_{n-1}) = \Gamma(q) + \Gamma(k)$. And, $\beta \Gamma(q) = \beta(a_0+4b_0, \dots, a_{n-1}+4b_{n-1}, b_0, \dots, b_{n-1}, c_0, \dots, c_{n-1}) = (\beta a_0 + 4\beta b_0, \dots, \beta a_{n-1} + 4\beta b_{n-1}, \beta b_0, \dots, \beta b_{n-1}, \beta c_0, \dots, \beta c_{n-1}) = \Gamma(\beta q)$.

Hence, Γ is a \mathbb{Z}_5 -linear map. Since Γ is a linear map, we have $\Gamma(q-k) = \Gamma(q) - \Gamma(k)$, for any $q, k \in S^n$. By the definition of the Lee distance, we have $d_L(q, k) = w_L(q-k) = w_H(\Gamma(q-k)) = w_H(\Gamma(q) - \Gamma(k)) = d_H(\Gamma(q), \Gamma(k))$. This shows that Γ is a distance preserving \mathbb{Z}_5 -linear map.

Theorem 2. *If C is a linear code of length n over S with cardinality $|C| = 5^k$ and Lee distance d_L , then the Gray image $\Gamma(C)$ is a $[5n, k, d_H]$ linear code over \mathbb{Z}_5 .*

Proof. The proof is given in article [7].

Example 1. $C_1 = \{0, u, 2u, 3u, 4u\}$ and $C_2 = \{0, v, 2v, 3v, 4v\}$ codes are linear codes of length 1 over S ring. Transforms $\Gamma(C_1)$ and $\Gamma(C_2)$ are linear codes $[5, 1, 2]$ and $[5, 1, 1]$ over \mathbb{Z}_5 , respectively.

Example 2. $C = \{(0, 0, 0, 0, 0), (v, v, v, v, v), (2v, 2v, 2v, 2v, 2v), (3v, 3v, 3v, 3v, 3v), (4v, 4v, 4v, 4v, 4v), (u, 0, 0, 0, 0), (2u, 0, 0, 0, 0), (3u, 0, 0, 0, 0), (4u, 0, 0, 0, 0), (u+v, v, v, v, v), (u+2v, 2v, 2v, 2v, 2v), (u+3v, 3v, 3v, 3v, 3v), (u+4v, 4v, 4v, 4v, 4v), (2u+v, v, v, v, v), (2u+2v, 2v, 2v, 2v, 2v), (2u+3v, 3v, 3v, 3v, 3v), (2u+4v, 4v, 4v, 4v, 4v), (3u+v, v, v, v, v), (3u+2v, 2v, 2v, 2v, 2v), (3u+3v, 3v, 3v, 3v, 3v), (3u+4v, 4v, 4v, 4v, 4v), (4u+v, v, v, v, v), (4u+2v, 2v, 2v, 2v, 2v), (4u+3v, 3v, 3v, 3v, 3v), (4u+4v, 4v, 4v, 4v, 4v)\}$

code is linear code of length 5 over S ring. Transform $\Gamma(C)$ is a $[25, 2, 2]$ linear code over \mathbb{Z}_5 .

Theorem 3. *Let Γ be the Gray map from S^n to \mathbb{Z}_5^{3n} . Let σ be the cyclic shift operator and ρ be the quasi-cyclic shift operator as defined in the preliminaries. Then $\Gamma\sigma = \rho\Gamma$.*

Proof. Let $q = (q_0, q_1, \dots, q_{n-1}) \in S^n$, where $q_i = a_i + ub_i + vc_i \in S$ and $a_i, b_i, c_i \in \mathbb{Z}_5$, for $i = 0, 1, \dots, n-1$.

Now $\Gamma(q) = (a_0+4b_0, a_1+4b_1, \dots, a_{n-1}+4b_{n-1}, b_0, b_1, \dots, b_{n-1}, c_0, c_1, \dots, c_{n-1})$. Applying ρ on both sides, we get

$\rho \Gamma(q) = \rho(a_0 + 4b_0, a_1 + 4b_1, \dots, a_{n-1} + 4b_{n-1}, b_0, b_1, \dots, b_{n-1}, c_0, c_1, \dots, c_{n-1})$
 $= (a_{n-1} + 4b_{n-1}, a_0 + 4b_0, \dots, a_{n-2} + 4b_{n-2}, b_{n-1}, b_0, \dots, b_{n-2}, c_{n-1}, c_0, \dots, c_{n-2}) \dots$
(1). On the other hand, we have $\Gamma \sigma(q) = \Gamma(q_{n-1}, q_0, \dots, q_{n-2}) = (a_{n-1} + 4b_{n-1}, a_0 + 4b_0, \dots, a_{n-2} + 4b_{n-2}, b_{n-1}, b_0, \dots, b_{n-2}, c_{n-1}, c_0, \dots, c_{n-2}) \dots$ (2).

Equality is obtained from (1) and (2).

Corollary 4. *Let C be a subset of S^n . Then C is a cyclic code of length n over S if and only if the Gray image $\Gamma(C)$ is a quasi-cyclic code of index 3 over \mathbb{Z}_5 with length $3n$.*

Proof. The proof is given in article [7].

Theorem 5. *Let Γ be the Gray map from S^n to \mathbb{Z}_5^{3n} , σ be the cyclic shift operator and Γ_π be the permutation version of the Gray map Γ as given before. Then $\Gamma_\pi \sigma = \sigma^3 \Gamma_\pi$.*

Proof. For any $q = (q_0, q_1, \dots, q_{n-1}) \in S^n$, where $q_i = a_i + ub_i + vc_i \in S$ and $a_i, b_i, c_i \in \mathbb{Z}_5$ for $i = 0, 1, \dots, n-1$. We have, $\sigma(q) = (q_{n-1}, q_0, \dots, q_{n-2})$. Applying Γ_π , we get $\Gamma_\pi \sigma(q) = \Gamma_\pi(q_{n-1}, q_0, q_1, \dots, q_{n-2}) = (\Gamma_\pi(q_{n-1}), \Gamma_\pi(q_0), \dots, \Gamma_\pi(q_{n-2})) = (a_{n-1} + 4b_{n-1}, b_{n-1}, c_{n-1}, a_0 + 4b_0, b_0, c_0, \dots, a_{n-2} + 4b_{n-2}, b_{n-2}, c_{n-2}) \dots$ (1)
On the other hand, we have $\Gamma_\pi(q) = (a_0 + 4b_0, b_0, c_0, a_1 + 4b_1, b_1, c_1, \dots, a_{n-1} + 4b_{n-1}, b_{n-1}, c_{n-1})$
 $\sigma \Gamma_\pi(q) = (c_{n-1}, a_0 + 4b_0, b_0, c_0, a_1 + 4b_1, b_1, c_1, \dots, a_{n-1} + 4b_{n-1}, b_{n-1})$
 $\sigma^2 \Gamma_\pi(q) = (b_{n-1}, c_{n-1}, a_0 + 4b_0, b_0, c_0, a_1 + 4b_1, b_1, c_1, \dots, a_{n-1} + 4b_{n-1})$
 $\sigma^3 \Gamma_\pi(q) = (a_{n-1} + 4b_{n-1}, b_{n-1}, c_{n-1}, a_0 + 4b_0, b_0, c_0, a_1 + 4b_1, b_1, c_1, \dots, a_{n-2} + 4b_{n-2}, b_{n-2}, c_{n-2}) \dots$
(2)

Equality is obtained from (1) and (2).

Corollary 6. *Let C be a subset of S^n . Then C is a cyclic code of length n over S if and only if $\Gamma_\pi(C)$ is equivalent to a 3-quasi-cyclic code of length $3n$ over \mathbb{Z}_5 .*

Proof. The proof is given in article [7].

4. CONSTACYCLIC CODES OVER S

Here, n -length λ -constacyclic codes on the S ring with $\lambda = (1 + 3u), (2 + 4u)$ and 4 unitary elements are examined. But in this part, $(1 + 3u)$ and $(2 + 4u)$ elements aren't provided transformations. Transform is provided for only 4 unitary elements.

Definition 4. *For $a \in \mathbb{Z}_5^{3n}$ with $a(a_0, a_1, \dots, a_{n-1}, a_n, \dots, a_{2n}, \dots, a_{3n-1}) = (a^{(0)}|a^{(1)}|a^{(2)})$, where $a^{(i)} \in \mathbb{Z}_5^n$ for $i = 0, 1, 2$, quasi-twisted shift operator on \mathbb{Z}_5^{3n} is defined by $v(a) = (\gamma_{(4)}(a^{(0)})|\gamma_{(4)}(a^{(1)})|\gamma_{(4)}(a^{(2)}))$, where $\gamma_{(4)}$ is a 4-constacyclic shift operator*

from \mathbb{Z}_5^n to \mathbb{Z}_5^n . A linear code C of length $3n$ over \mathbb{Z}_5 is called a quasi-twisted code of index 3 if $v(C) = C$.

Theorem 7. Let $\gamma_{(4)}$ be 4-constacyclic shift operator, Γ be the Gray map and v be the quasi-twisted shift operator as given before. Then $\Gamma \gamma_{(4)} = v \Gamma$.

Proof. Let $q = (q_0, q_1, \dots, q_{n-1}) \in S^n$, where $q_i = a_i + ub_i + vc_i \in S$ and $a_i, b_i, c_i \in \mathbb{Z}_5$, for $i = 0, 1, \dots, n-1$. Then $\Gamma \gamma_{(4)}(q) = \Gamma(4q_{n-1}, q_0, \dots, q_{n-2}) = \Gamma(4a_{n-1} + u(4b_{n-1}) + v(4c_{n-1}), a_0 + ub_0 + vc_0, \dots, a_{n-2} + ub_{n-2} + vc_{n-2}) = (4a_{n-1} + b_{n-1}, a_0 + 4b_0, \dots, a_{n-2} + 4b_{n-2}, 4b_{n-1}, b_0, \dots, b_{n-2}, 4c_{n-1}, c_0, \dots, c_{n-2}) \dots$ (1) On the other hand, we have $v \Gamma(q) = v(a_0 + 4b_0, a_1 + 4b_1, \dots, a_{n-1} + 4b_{n-1}, b_0, b_1, \dots, b_{n-1}, c_0, c_1, \dots, c_{n-1}) = (4(a_{n-1} + 4b_{n-1}), a_0 + 4b_0, \dots, a_{n-2} + 4b_{n-2}, 4b_{n-1}, b_0, \dots, b_{n-2}, 4c_{n-1}, c_0, \dots, c_{n-2}) = (4a_{n-1} + b_{n-1}, a_0 + 4b_0, \dots, a_{n-2} + 4b_{n-2}, 4b_{n-1}, b_0, \dots, b_{n-2}, 4c_{n-1}, c_0, \dots, c_{n-2}) \dots$ (2) Equality is obtained from (1) and (2).

Corollary 8. A code C is a 4-constacyclic code over S if and only if $\Gamma(C)$ is a quasi-twisted code of index 3 over \mathbb{Z}_5 with length $3n$.

Proof. The proof is given in article [7].

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