

## Exercise Sheet 6

The main theme of this sheet is to introduce the so-called *Zig-Zag product*, which will give us an explicit combinatorial construction of expander graphs. For the whole sheet,  $G$  will be a “large”  $D$ -regular graph on  $N$  vertices, and  $H$  will be a “small”  $d$ -regular graph on  $D$  vertices. Furthermore, we assume that the neighbors of each vertex of  $G$  have a prescribed order. So for instance, if a vertex  $v \in V(G)$  has neighbors  $u_1, \dots, u_D$ , we will use the  $i$ -th neighbor of  $v$  to denote the vertex  $u_i$ . Note that if  $u$  and  $v$  are two adjacent vertices in  $G$ , it might be that  $u$  is the  $i$ -th neighbor of  $v$ , and  $v$  is the  $j$ -th neighbor of  $u$  for  $i \neq j$ . We will also assume that the vertex-set of  $H$  is  $\{1, 2, \dots, D\}$ .

1. Before introducing the Zig-Zag product itself, let us start with a simpler notion of a product of graphs  $G$  and  $H$  called the *replacement product*. We define the result of the replacement product  $G \boxtimes H$  to be the following  $(d+1)$ -regular graph with  $N \cdot D$  vertices:

- The vertex set will be  $V(G) \times V(H)$ .
- For any vertex  $u \in V(G)$ , the vertices  $(u, i)$  and  $(u, j)$  of  $G \boxtimes H$  where  $\{i, j\} \subseteq V(H)$ , will form an edge if and only if  $\{i, j\}$  is an edge of  $H$ .
- Finally, two vertices  $(u, i)$  and  $(v, j)$  for  $u \neq v$  will be connected by an edge if and only if  $\{u, v\}$  is an edge of  $G$ ,  $v$  is the  $i$ -th neighbor of  $u$ , and  $u$  is the  $j$ -th neighbor of  $v$ .

In other words, we expand every vertex of  $V(G)$  into  $D$  vertices (each such expansion is called *cloud*), each original edge  $\{u, v\}$  of  $G$  will become an edge between  $(u, i)$  and  $(v, j)$ , where  $i$  is the integer denoting the position of  $v$  among neighbors of  $u$ , analogously  $u$  is the  $j$ -th neighbor of  $v$ , and finally inside each cloud we place a copy of  $H$ .

- a) Determine the graph  $K_{d+1} \boxtimes K_d$ .
- b) Show that if  $G$  is connected and  $H$  is connected, then  $G \boxtimes H$  will also be a connected graph.

2. Now we introduce an analogous product that keeps not only the connectivity, but also good expansion. The product is called Zig-Zag product, denoted by  $G \circledast H$ . The vertex-set of  $G \circledast H$  will be again  $V(G) \times V(H)$ , and the edges will correspond to the following paths of length 3 in  $G \oplus H$ . Two vertices  $(u, i)$  and  $(v, \ell)$  of  $G \circledast H$  will be connected by an edge if and only if there exist  $j \in V(H)$  and  $k \in V(H)$  such that:

- $(u, i)$  is connected to  $(u, j)$  in  $G \oplus H$  (i.e.,  $\{i, j\}$  is an edge of  $H$ ),
- $(u, j)$  is connected to  $(v, k)$  in  $G \oplus H$  (i.e.,  $v$  is the  $j$ -th neighbor of  $u$ , and  $u$  is the  $k$ -th neighbor of  $v$  in  $G$ ), and
- $(v, k)$  is connected to  $(v, \ell)$  in  $G \oplus H$  (i.e.,  $\{k, \ell\}$  is an edge of  $H$ ).

Loosely speaking, the edges of  $G \circledast H$  exactly correspond to “zig-zag” paths between two clouds in  $G \oplus H$ .

- a) Show that the  $G \circledast H$  is a  $d^2$ -regular graph.
- b) What is the graph  $K_{d+1} \circledast K_d$ ?
- c) Again, if  $G$  is connected and  $H$  is connected, show that  $G \circledast H$  is also connected.

3. Ok, we introduced the required notion, now it is time for a serious stuff. Prove that if  $\lambda_{\max}(G)$ , which is the second largest eigenvalue in absolute value of the adjacency matrix of  $G$  (i.e.,  $\lambda_{\max} = \max\{\lambda_2(A(G)), -\lambda_N(A(G))\}$  for  $A(G)$  the adjacency matrix of  $G$ ) is equal to  $\alpha \cdot D$ , and analogously  $\lambda_{\max}(H) = \beta \cdot d$ , then

$$\lambda_{\max}(G \circledast H) \leq d^2 \cdot (\alpha + \beta + \beta^2) .$$

*Hint: Decompose the adjacency matrix  $A$  of  $G \circledast H$  into three pieces using the interpretation of the edges in  $G \circledast H$  as paths  $G \oplus H$  which go inside a first cloud following edges of  $H$ , then between two clouds following edges of  $G$ , and finally inside the second cloud again following edges of  $H$ . Then recall the Rayleigh formula and the spectral norm of a matrix, and use them to bound the  $\lambda_{\max}(G \circledast H)$ .*

4. So if  $G$  and  $H$  are good expanders, we showed that the Zig-Zag product outputs also an expander. The final step is to appropriately use the product and construct a sequence of bounded-degree expanders. Suppose that  $H$  is a  $d$ -regular graph on  $d^4$  vertices with  $\lambda_{\max}(H) = d/4$  (such a graph exists, either use the probabilistic method to show that, or take it as granted; if needed, you could brute-force all  $d$ -regular graphs with  $d^4$  vertices in order to find this one graph).

Let  $G_0 := H$ ,  $G_1 := (G_0)^2$  and for  $k \geq 2$ , let  $G_k := (G_{k-1})^2 \circledast H$ . The square of a graph  $G$ , denoted by  $G^2$ , is the graph on the same set of vertices as  $G$ , where the edges correspond to walks of length two in  $G$  (so in particular, if  $G$  is a  $d$ -regular graph without loops,  $G^2$  is a  $d^2$ -regular graph containing exactly  $d$  loops on each vertex).

Show that for all  $k \geq 1$ ,  $G_k$  is  $d^2$ -regular graph with  $d^{4k}$  vertices and  $\lambda_{\max}(G_k) \leq d^2/2$ .

**Submission: Wednesday, 20th April 2016 during the exercise class.**