

# Graph Theory

## Solutions 1

*The aim of the homework problems is to help you understand the theory better by actively using it to solve exercises. **Do not read the solutions** before you believe you have solved the problems: it ruins your best way of preparing for the exam. The purpose of this write-up is merely to provide some guideline on how solutions should look like, and to help clean up hazy arguments. For hints, feel free to consult your teaching assistant.*

**Problem 1(c):** Let us call the vertices  $a, b, c, d, e, f, g, h$  and suppose they are ordered by their degrees (so  $d(a) = 1$ ,  $d(b) = 2$  and so on). We have 8 vertices, so each one of degree 6 is connected to all but one of the other vertices. So each of  $e, f, g$  and  $h$  sends at least one edge to the set  $\{a, b\}$ , but that is impossible, since  $a$  and  $b$  are adjacent to at most 3 edges in total. Hence there is no such graph.

**Problem 5:** We do induction on  $n$ . If  $n = 7$  then we have  $35 - 14 = 21$  edges. The only graph on 7 vertices with 21 edges is  $K_7$ , so we can take  $K_7$  as a subgraph satisfying the conditions.

Now suppose that  $n > 7$  and we know the statement for  $n - 1$ . If all the degrees in the graph are at least 6, we can again take the graph itself is a good subgraph. If not, then there is a vertex  $v$  with  $d(v) \leq 5$ . Then  $G - v$  is a graph on  $n - 1$  vertices, and it has at least  $5n - 14 - 5 = 5(n - 1) - 14$  edges (we lose at most 5 by deleting  $v$ ). So by induction  $G - v$  has a good subgraph, but that is also a good subgraph of  $G$ .

**Problem 6:** Suppose there are two vertex-disjoint paths  $P_1$  and  $P_2$  of maximum length  $l$ . Since the graph is connected, there is a path  $Q$  between  $P_1$  and  $P_2$ , say from  $w_1 \in P_1$  to  $w_2 \in P_2$  such that the interior vertices of  $Q$  avoid  $P_1$  and  $P_2$  (or formally:  $V(Q) \cap (V(P_1) \cup V(P_2)) = \{w_1, w_2\}$ ). Here  $w_1$  and  $w_2$  cut  $P_1$  and  $P_2$  into two pieces each. Let  $P'_1$  be the longer piece of  $P_1$  and  $P'_2$  be the longer piece of  $P_2$ , so both  $P'_1$  and  $P'_2$  have length at least  $l/2$ . Moreover,  $P'_1 \cup Q \cup P'_2$  forms a path(!) of length at least  $l + e(Q)$ , which is impossible because the longest path had length  $l$ . This is a contradiction.

To say that  $P'_1 \cup Q \cup P'_2$  is a path, we really needed the fact that  $Q$  is internally disjoint from  $P_1$  and  $P_2$ . But why is there such a  $Q$ ? To show this, take arbitrary vertices  $v_1 \in P_1$  and  $v_2 \in P_2$ . As  $G$  is connected, there is a path  $Q_0$  from  $v_1$  to  $v_2$ . Since  $P_1$  and  $P_2$  are vertex-disjoint, there is a *last* vertex in  $Q_0$  from  $P_1$ , let us call this  $w_1$ . We define  $w_2$  to be the first vertex from

$P_2$  appearing after  $w_1$  in  $Q_0$ . Then the part of the path  $Q_0$  between  $w_1$  and  $w_2$  is a good choice for  $Q$ .