

HOMOTOPICAL AND HIGHER ALGEBRA - SEMINAR

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Week 4 - Extending down TFTs

M closed oriented n -mfd. Consider it as $M \in \text{Mor}(\emptyset, \emptyset)$

$$\begin{array}{c} \downarrow Z \\ Z(M): Z(\emptyset) \xrightarrow{\cong k} Z(\emptyset) \in \text{Hom}(k, k) \cong k \end{array}$$

$\hookrightarrow Z$ assigns to M an ~~non~~ element of k , a diffeo. invariant.

NEW POINT OF VIEW: these diffeo. invariants are our new main object of interest.

Since we can also evaluate Z on n -mfd with boundary and on $(n-1)$ -mfd without boundary, we try to compute $Z(M)$ by breaking M into pieces.

Ex. Σ_g : closed oriented 2-mfd of genus $g \geq 0$

Recall $\text{Cat}(\text{comm. Frobenius algebras}) \cong \text{2TFT}$

Given: a Frob. alg A with basis $\{T_0, T_1, \dots, T_r\}$ such that $T_0 = 1_A$

multiplication $\mu: A \otimes A \rightarrow A$

defined by $\mu(T_i \otimes T_j) = \mu_{ij}^k T_k$

Frobenius form $\epsilon: A \rightarrow k$

We define an associative nondegenerate pairing $\beta: A \otimes A \rightarrow k$

$(x \otimes y) \mapsto \epsilon(\mu(x \otimes y))$

by $\beta(T_i \otimes T_j) = \beta_{ij}$

nondegeneracy of $\beta \Rightarrow \exists \gamma: k \rightarrow A \otimes A$

defined by $\gamma(1) = \gamma^{ij} T_i \otimes T_j$

such that $\beta_{ij} \gamma^{jk} = \delta_i^k$

$\gamma^{ij} \beta_{jk} = \delta_k^i$ Kronecker delta

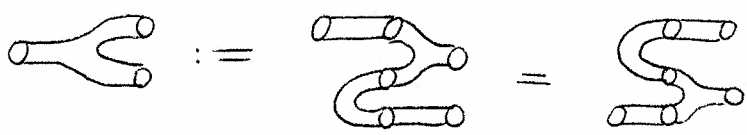
i.e. $(\beta_{ij})_{i,j=1}^n \text{ inverse } (\gamma^{ij})_{i,j=1}^n$

We've seen that given a Froben algebra we obtain a ZTFT by setting

$$Z(S^1) = A, \quad Z(\text{cup}) = \mu, \quad Z(\text{cap}) = \varepsilon, \quad Z(\text{box}) = \beta, \quad Z(\text{circle}) = \gamma$$

and $Z(\text{point}) = k \rightarrow A$
 $1_k \mapsto 1_A$

We also want to define a comultiplication $\delta: A \rightarrow A \otimes A$ as



Hence we define $\delta(T_k) = \delta_k^{ij} T_i \otimes T_j$
 by $\delta_k^{ij} = \gamma^{ie} \mu_{ek}^d = \mu_{kf}^i \gamma^{fj}$
 and set $Z(\text{circle}) = \delta$.

CONCRETE EX. \mathbb{C} as vect. space over \mathbb{R} with basis $T_0=1, T_1=i$

usual multiplication
 $\varepsilon: \mathbb{C} \rightarrow \mathbb{R}$
 $z \mapsto \text{Re}(z)$
 \updownarrow
 $\beta: \mathbb{C} \otimes \mathbb{C} \rightarrow \mathbb{R}$
 $z_1 \otimes z_2 \mapsto x_1 x_2 - y_1 y_2$
 ($\text{Re}(zi) = x_i, \text{Im}(zi) = y_i$)

computations... \Rightarrow

$$\mu_{00}^0 = 1, \mu_{00}^1 = 0, \mu_{01}^0 = 0, \mu_{01}^1 = 1$$

$$\mu_{10}^0 = 0, \mu_{10}^1 = 1, \mu_{11}^0 = -1, \mu_{11}^1 = 0$$

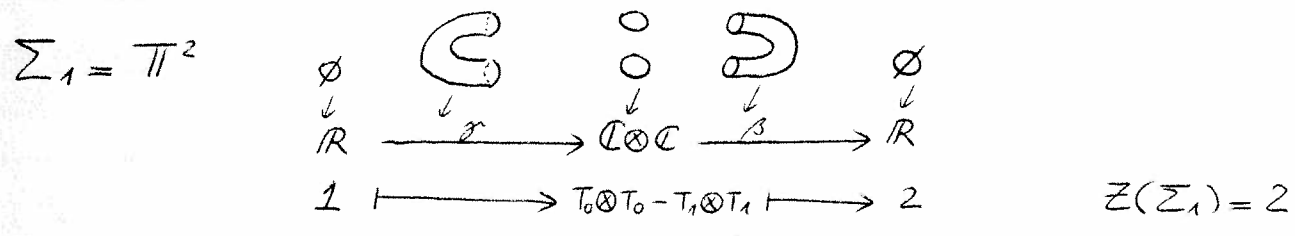
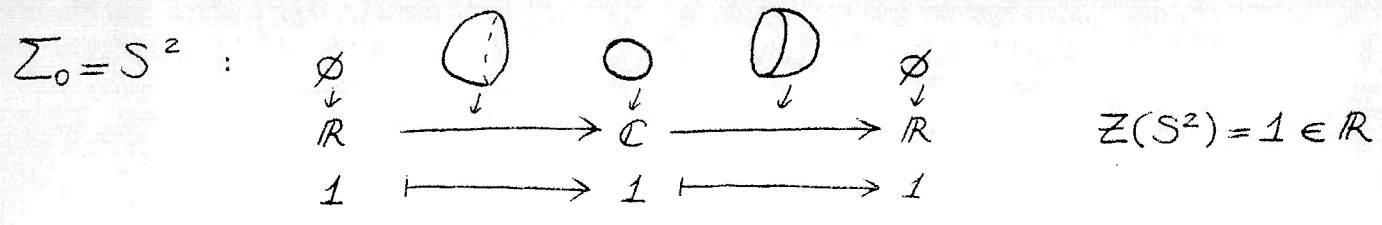
$$\beta_{00} = 1, \beta_{10} = 0, \beta_{01} = 0, \beta_{11} = -1$$

$$(\gamma) = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

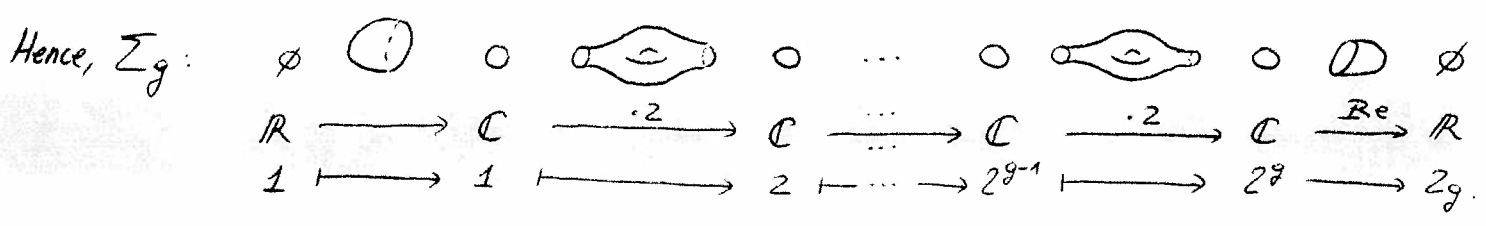
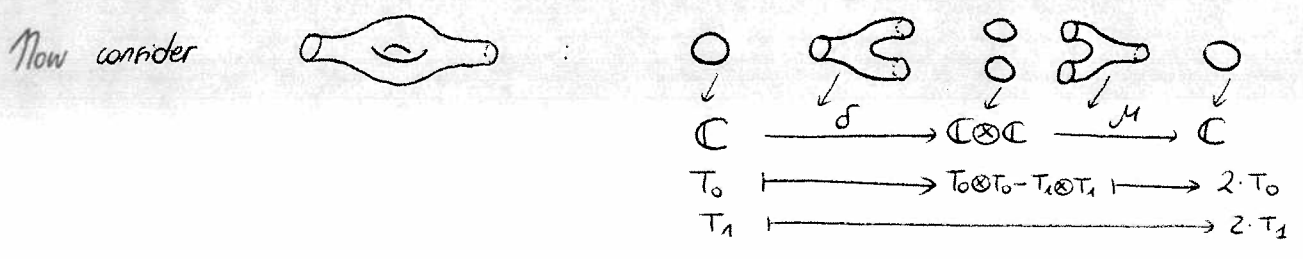
$$\delta: \mathbb{C} \otimes \mathbb{C} \rightarrow \mathbb{C} \otimes \mathbb{C}$$

$$T_0 \mapsto T_0 \otimes T_0 - T_1 \otimes T_1$$

$$T_1 \mapsto T_0 \otimes T_1 + T_1 \otimes T_0$$



(in fact for any Froben. alg $1 \xrightarrow{\alpha} \gamma_{ij} T_i \otimes T_j \xrightarrow{\beta} \beta_{ij} \gamma_{ij} = n = \dim(A)$)



$\Rightarrow Z(\Sigma_g) = 2^g \in \mathbb{R}$

We try to go on with this process for $n > 2$.

M oriented n -mfd, consider it as $M \in \text{Mor}(\emptyset, \partial M)$

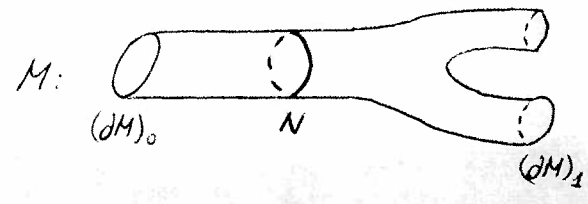
$$\downarrow z$$

$$Z(M): Z(\emptyset) \longrightarrow Z(\partial M)$$

$\cong k$

we interpret it as $Z(M) \in Z(\partial M)$
 \uparrow vector space.

In order to translate functoriality:



take a closed submfd N of codim 1

$$M_0: \emptyset \quad \begin{array}{c} \textcircled{\partial M}_0 \\ \textcircled{N} \\ \textcircled{\partial M}_1 \end{array}$$

$$k \longrightarrow Z((\partial M)_0) \otimes Z(N)$$

$$\downarrow$$

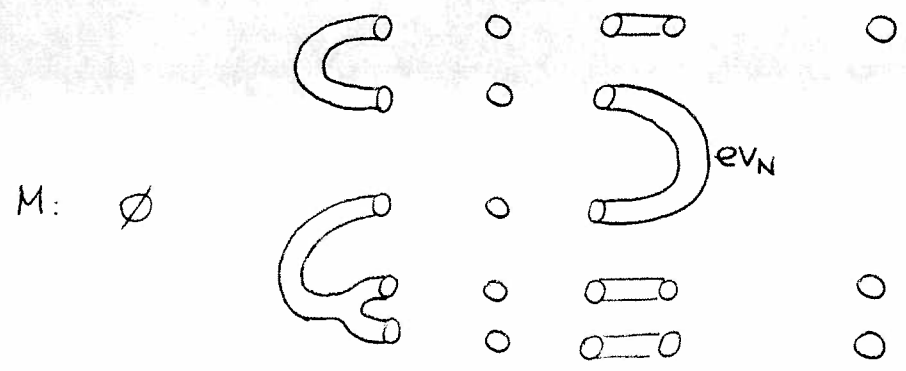
$$Z(M_0)$$

$$M_1: \emptyset \quad \begin{array}{c} \textcircled{\bar{N}} \\ \textcircled{\partial M}_1 \end{array}$$

$$k \longrightarrow Z(\bar{N}) \otimes Z((\partial M)_1)$$

$$\downarrow$$

$$Z(M_1)$$



$$k \longrightarrow Z((\partial M)_0) \otimes Z(N) \otimes Z(\bar{N}) \otimes Z((\partial M)_1) \longrightarrow Z(\partial M)$$

$$\cong$$

$$Z(\partial M) \otimes Z(N) \otimes Z(\bar{N})$$

$$1 \longmapsto Z(M_0) \otimes Z(M_1) \longmapsto Z(M)$$

\Rightarrow We have found a rule for computing $Z(M)$ given a decomposition $M = M_0 \amalg_N M_1$
 $\# N$ is closed and of codim 1.

We've seen that for $n=2$ is enough (only needed pieces $\square, \bigcirc, \bigtriangleright, \mathbb{R}^2, \mathbb{S}^1$)

Can we find $\forall n$ a list of "simple" n -mfld with boundary such that any n -mfld is obtained by gluing these?

For $n \geq 3$ this gets complicated:

- we don't have a complete classification of n -mflds
- $\# n \gg 0$, ~~this~~ ^{this process} doesn't help us a lot, since we don't know how the submflds of codim 1 look like.

We could use triangulations, but the gluing gets more complicated (we have "corners")

\rightarrow We will want to allow us to glue along submanifolds which have boundary themselves.

FIRST ATTEMPT TO EXTEND
 THE DEFINITION OF A TFT:

We need these data:

(1) M^n closed oriented $\Rightarrow Z(M) \in k$

(2) N^{n-1} closed oriented \Rightarrow a k -vect. space $Z(N)$, if $N = \emptyset$: $Z(N) = Z(\emptyset) \cong k$ canonically

(3) M^n oriented \Rightarrow an element $Z(M) \in Z(\partial M)$, if $\partial M = \emptyset$ $Z(M)$ should correspond to the element given in (1) under the \cong from (2).

(4) P^{n-2} closed oriented $\Rightarrow Z(P)$ a k -linear category, i.e. $Z(P)$ a category st. $\forall X, Y \in \text{Obj}(Z(P))$
 $\text{Mor}_{Z(P)}(X, Y)$ is a k -vect. space
 and
 $\text{Mor}_{Z(P)}(X, Y) \times \text{Mor}_{Z(P)}(Y, Z) \rightarrow \text{Mor}_{Z(P)}(X, Z)$
 is bilinear
 , if $P = \emptyset$ then $Z(P) = Z(\emptyset) = \text{Vect}(k)$.

(5) N^{n-1} oriented \Rightarrow an object $Z(N) \in \text{Obj}(Z(\partial N))$, if $\partial N = \emptyset$: $Z(N)$ should correspond to the vect. space $Z(N) \in \text{Obj}(\text{Vect}(k))$ given in (2).

This first attempt is not enough we have some

PROBLEMS: • $\dim(B) = n$, B a bordism

we didn't specify a method for defining a linear map $Z(B): Z(M) \rightarrow Z(N) \neq M \neq \emptyset$

• $\dim(B) = n-1$, B a bordism

we didn't specify how to define the functor $Z(B): Z(M) \rightarrow Z(N)$ when $M \neq \emptyset$.

• we don't have any condition for gluings, nor for disjoint unions.

We try to refine the definition.

A STRICT 2-CATEGORY \mathcal{C} consists of:

• a collection of objects X, Y, Z, \dots

• $\forall X, Y$ objects, $\text{Map}_{\mathcal{C}}(X, Y)$ is a category

• $\forall X$ object, a distinct object $\text{id}_X \in \text{Map}_{\mathcal{C}}(X, X)$

• $\forall X, Y, Z$ objects, a composition functor

$$\Psi: \text{Map}_{\mathcal{C}}(X, Y) \times \text{Map}_{\mathcal{C}}(Y, Z) \longrightarrow \text{Map}_{\mathcal{C}}(X, Z)$$

$$\Psi(F \circ F', G \circ G') = \Psi(F, G) \circ \Psi(F', G')$$

• $\Psi(F, \text{id}_Y) = F$

$$\Psi(\text{id}_X, G) = G$$

• $\text{Map}_{\mathcal{C}}(W, X) \times \text{Map}_{\mathcal{C}}(X, Y) \times \text{Map}_{\mathcal{C}}(Y, Z) \longrightarrow \text{Map}_{\mathcal{C}}(W, Y) \times \text{Map}_{\mathcal{C}}(Y, Z)$

$$\begin{array}{ccc} \downarrow & \curvearrowright & \downarrow \\ \text{Map}_{\mathcal{C}}(W, X) \times \text{Map}_{\mathcal{C}}(X, Z) & \longrightarrow & \text{Map}_{\mathcal{C}}(W, Z) \end{array}$$

This is also called a CATEGORY ENRICHED IN CATEGORIES.

this means that $\forall X, Y$ objects
 $\text{Map}_{\mathcal{C}}(X, Y)$ is a ~~category~~
category and satisfy the
listed conditions

Remark In fact we can enrich a category \mathcal{C} in any monoidal category \mathcal{M} by requiring that $\text{Map}_{\mathcal{C}}(X, Y)$ is an object in \mathcal{M} . For example: ordinary categories are enriched in **Sets**, k -linear categories are enriched in $\text{Vect}(k)$.

Example $\text{Vect}_2(k)$

• OBJECTS cocomplete k -linear categories $\mathcal{A}, \mathcal{B}, \mathcal{C}, \dots$

i.e. \mathcal{A} with objects X, Y, Z, \dots is such that

- k -linear $\left\{ \begin{array}{l} \bullet \text{Hom}_{\mathcal{A}}(X, Y) \text{ is } k\text{-vect. space } \forall X, Y \\ \bullet \text{Hom}_{\mathcal{A}}(X, Y) \times \text{Hom}_{\mathcal{A}}(Y, Z) \longrightarrow \text{Hom}_{\mathcal{A}}(X, Z) \text{ is bilinear} \end{array} \right.$
 - cocomplete $\left\{ \begin{array}{l} \bullet \text{closed under } \oplus \\ \bullet \text{closed under } \text{coker}(-) \end{array} \right.$
- \nearrow
all colimits exist

• $\text{Map}_{\text{Vect}_2(k)}(\mathcal{A}, \mathcal{B})$ is the category of cocomplete k -linear ~~functors~~ functors
 i.e. ~~functors~~ ~~functors~~ ~~with~~ ~~with~~
 (cocomplete \iff commutes with colimits)

- $\text{id}_{\mathcal{A}}$ the identity functor
- composition clear

Example $\text{Cob}_2(n)$ NOT A STRICT 2-CATEGORY

we try to define it:

• OBJECTS closed oriented $(n-2)$ -manifolds M, N, P, \dots

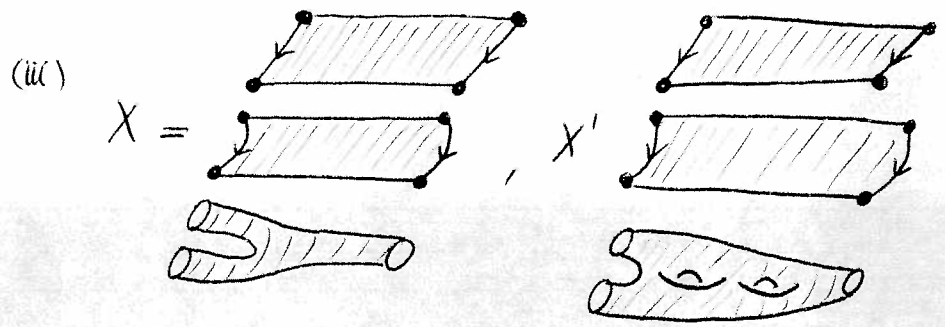
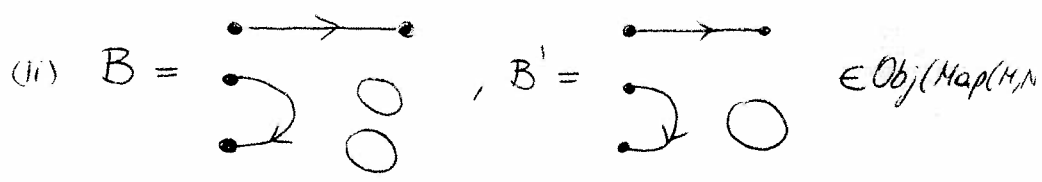
• $\text{Map}_{\text{Cob}_2(n)}(M, N)$ category with - OBJECTS bordisms B from M to N

i.e. B^{n-1} mfd together with diffeo $\partial B \cong \overline{M} \sqcup N$

- MORPHISMS $\forall B, B' \in \text{Obj}(\text{Map}_{\text{Cob}_2(n)}(M, N))$

$\text{Hom}(B, B') = \{ \text{oriented diffeo. classes of oriented bordisms from } B \text{ to } B' \text{ such that the diffeos are trivial on the common boundary } \partial B \cong \overline{M} \sqcup N = \partial B' \}$

Picture/Ex. (i) $M = \begin{matrix} + \cdot \\ + \cdot \\ - \cdot \end{matrix}$ $N = \cdot +$



$X, X' \in \text{Hom}(B, B')$, $X \neq X'$

Rem. $\partial X = \partial X' = \overline{B'} \amalg_{M \amalg N} ((M \amalg N) \times [0, 1]) \amalg_{M \amalg N} B'$

As pointed out $\text{Cob}_2(n)$ is not a strict 2-category. We encounter some difficulties:

- (1) given $B: M \rightarrow M'$ and $B': M' \rightarrow M''$ we want to define a composition by gluing along M' . We need a well defined composition, not up to diffeo. \Rightarrow we need to define a smooth structure on the composition (for ex. fix a smooth collar nbhd of M' in B and B').

(2) given B, B' and $B'': M'' \rightarrow M'''$ we want

$$(B \amalg_{M'} B') \amalg_{M''} B'' \stackrel{!}{=} B \amalg_{M'} (B' \amalg_{M''} B'')$$

\uparrow we can arrange to have a diffeo, but equality is difficult.

We will change this definition in order to fix these problems.

For the moment, suppose we have a good definition that takes care of these problems.

Def. A 2-EXTENDED TOPOLOGICAL FIELD THEORY of dim. n is a symmetric monoidal functor $Z: \text{Cob}_2(n) \rightarrow \text{Vect}_2(k)$ between 2-categories.

\hookrightarrow we need to understand $\text{Cob}_2(n)$ and $\text{Vect}_2(k)$ as symmetric monoidal 2-categories.

For $\text{Cob}_2(n)$ it's easy $\otimes := \amalg$
 $\mathbb{1}_{\text{Cob}_2(n)} = \emptyset^{(n-2)}$

For $\text{Vect}_2(k)$ it's more complicated. We can define a \otimes between categories and then $\mathbb{1}_{\text{Vect}_2(k)} = \text{Vect}(k)$

We called this new functor extended, because in fact it incorporates our old definition of nTFT:

For \mathcal{C} a symm. mon. 2-category

$\Rightarrow \Omega \mathcal{C} := \text{Map}_{\mathcal{C}}(\mathbb{1}_{\mathcal{C}}, \mathbb{1}_{\mathcal{C}})$ is a symm. mon. category

We see that $\Omega \text{Cob}_2(n) = \text{Map}_{\text{Cob}_2(n)}(\emptyset^{n-2}, \emptyset^{n-2}) \cong \text{Cob}(n)$

$\Omega \text{Vect}_2(k) = \text{Map}_{\text{Vect}_2(k)}(\text{Vect}(k), \text{Vect}(k)) \cong \text{Vect}(k)$.

Notice that ΩZ allows us to make computations by cutting along closed submfld of codim 1, Z , on the other hand, allows us to cut along submflds of codim 1 with boundary.

We would like to go on with this process, and cut N along a submfld P , $N = N_0 \amalg_P N_1$, but this makes sense in our definition only if P is closed.

So, the idea is to extend TFT's even more. We start with a definition.

Def. For $n \geq 0$, a STRICT n -CATEGORY is defined by induction on n :

$n=0$ a strict 0-category is a set

$n > 0$ a strict n -category is a category enriched over strict $(n-1)$ -categories

i.e. we have the following

- ① a collection of objects X, Y, Z, \dots
- ② $\forall X, Y$ objects
 $\text{Map}(X, Y)$ is a strict $(n-1)$ -category
- ③ $\text{id}_X \in \text{Map}(X, X)$ identity objects
 $\text{Map}(X, Y) \times \text{Map}(Y, Z) \rightarrow \text{Map}(X, Z)$ composition
 satisfying unity and associativity conditions.

- Notice
- $n=1$ we have a category enriched over set, i.e. a usual category
 - $n=2$ the definition of strict 2-category given above.

We will weaken this definition in order to have associativity up to isomorphism for the composition, and not strict equality.